

An efficient parallel algorithm  
for estimating  
*higher-order polyspectra*

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PTChat@Kyoto, 10 April 2019

# Work done by



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# Higher order polyspectra

- In an statistically homogeneous universe

- Bispectrum

$$\langle \delta(\mathbf{k}_1)\delta(\mathbf{k}_2)\delta(\mathbf{k}_3) \rangle = (2\pi)^3 B(\mathbf{k}_1, \mathbf{k}_2, \mathbf{k}_3) \delta^D(\mathbf{k}_1 + \mathbf{k}_2 + \mathbf{k}_3)$$

- Trispectrum

$$\langle \delta(\mathbf{k}_1)\delta(\mathbf{k}_2)\delta(\mathbf{k}_3)\delta(\mathbf{k}_4) \rangle = (2\pi)^3 T(\mathbf{k}_1, \mathbf{k}_2, \mathbf{k}_3, \mathbf{k}_4) \delta^D(\mathbf{k}_1 + \mathbf{k}_2 + \mathbf{k}_3 + \mathbf{k}_4)$$

- Quadspectrum, Pentaspectrum, etc

$$\langle \delta(\mathbf{k}_1)\delta(\mathbf{k}_2) \cdots \delta(\mathbf{k}_n) \rangle = (2\pi)^3 P_n(\mathbf{k}_1, \mathbf{k}_2, \cdots, \mathbf{k}_n) \delta^D(\mathbf{k}_1 + \mathbf{k}_2 + \cdots + \mathbf{k}_n)$$

# Naive estimator: bispectrum

- Let's focus on the monopole, because the extension is trivial
  - From the definition:

$$\langle \delta(\mathbf{k}_1)\delta(\mathbf{k}_2)\delta(\mathbf{k}_3) \rangle = (2\pi)^3 B(\mathbf{k}_1, \mathbf{k}_2, \mathbf{k}_3) \delta^D(\mathbf{k}_1 + \mathbf{k}_2 + \mathbf{k}_3)$$

- We can estimate the bispectrum from direct sampling:

$$B(k_1, k_2, k_3) = \frac{V_f}{V_{(123)}^B (2\pi)^3} \int_{k_1} d^3 q_1 \int_{k_2} d^3 q_2 \delta(\mathbf{q}_1) \delta(\mathbf{q}_2) \delta(-\mathbf{q}_{12})$$

# Complexity of Naive estimator

- To measure bispectrum from  $(k_1, k_2, k_3) = (1, 1, 1)k_F$  to  $(k_1, k_2, k_3) = (N_{\max}, N_{\max}, N_{\max})k_F$ , we need to loop over  $k_{1x}, k_{1y}, k_{1z}, k_{2x}, k_{2y}, k_{2z}$  (then,  $k_3$  is determined from triangle condition)
- Complexity =  $(N_{\max})^6$

$$B(k_1, k_2, k_3) = \frac{V_f}{V_{(123)}^B (2\pi)^3} \int_{k_1} d^3 q_1 \int_{k_2} d^3 q_2 \delta(\mathbf{q}_1) \delta(\mathbf{q}_2) \delta(-\mathbf{q}_{12})$$

# Roman's estimator

$$\begin{aligned} B(k_1, k_2, k_3) &= \frac{V_f}{V_{(123)}^B (2\pi)^3} \int_{k_1} d^3 q_1 \int_{k_2} d^3 q_2 \int_{k_3} d^3 q_3 \delta(\mathbf{q}_1) \delta(\mathbf{q}_2) \delta(\mathbf{q}_3) \delta_D(\mathbf{q}_{123}) \\ &= \frac{V_f}{V_{(123)}^B (2\pi)^3} \int_{k_1} d^3 q_1 \int_{k_2} d^3 q_2 \int_{k_3} d^3 q_3 \delta(\mathbf{q}_1) \delta(\mathbf{q}_2) \delta(\mathbf{q}_3) \int \frac{d^3 x}{(2\pi)^3} e^{i\mathbf{x} \cdot \mathbf{q}_{123}} \\ &= \frac{V_f}{V_{(123)}^B (2\pi)^3} \int \frac{d^3 x}{(2\pi)^3} \left( \int_{k_1} d^3 q_1 \delta(\mathbf{q}_1) e^{i\mathbf{x} \cdot \mathbf{q}_1} \right) \left( \int_{k_2} d^3 q_2 \delta(\mathbf{q}_2) e^{i\mathbf{x} \cdot \mathbf{q}_2} \right) \left( \int_{k_3} d^3 q_3 \delta(\mathbf{q}_3) e^{i\mathbf{x} \cdot \mathbf{q}_3} \right) \\ &= \frac{V_f}{V_{(123)}^B (2\pi)^3} \int \frac{d^3 x}{(2\pi)^3} I_{k_1}(\mathbf{x}) I_{k_2}(\mathbf{x}) I_{k_3}(\mathbf{x}) \end{aligned}$$

$$I_{k_i}(\mathbf{x}) = \int_{k_i} d^3 q_1 \delta(\mathbf{q}) e^{i\mathbf{x} \cdot \mathbf{q}} = \int d^3 q_1 \tilde{I}_{k_i}(\mathbf{q}) e^{i\mathbf{x} \cdot \mathbf{q}}$$

# Complexity of Roman's estimator

- To measure bispectrum from  $(k_1, k_2, k_3) = (1, 1, 1)k_F$  to  $(k_1, k_2, k_3) = (N_{\max}, N_{\max}, N_{\max})k_F$ , we need to loop over  $k_1, k_2, k_3$  and need to calculate the inner product of 3D matrices
- Complexity =  $(N_{\max})^6$

$$B(k_1, k_2, k_3) = \frac{V_f}{V_{(123)}^B (2\pi)^3} \int \frac{d^3 x}{(2\pi)^3} I_{k_1}(\mathbf{x}) I_{k_2}(\mathbf{x}) I_{k_3}(\mathbf{x})$$

# Yet, Roman's estimator is much faster!

- Loop is only for  $k_1, k_2, k_3$  : much fewer computation
- Matrix inner product is much faster than irregular sampling of the matrix

$$B(k_1, k_2, k_3) = \frac{V_f}{V_{(123)}^B (2\pi)^3} \int_{k_1} d^3 q_1 \int_{k_2} d^3 q_2 \delta(\mathbf{q}_1) \delta(\mathbf{q}_2) \delta(-\mathbf{q}_{12})$$

$$B(k_1, k_2, k_3) = \frac{V_f}{V_{(123)}^B (2\pi)^3} \int \frac{d^3 x}{(2\pi)^3} I_{k_1}(\mathbf{x}) I_{k_2}(\mathbf{x}) I_{k_3}(\mathbf{x})$$



# Memory requirement

- To get an unbiased bispectrum,  
 $N_{\text{mesh}} > 3 N_{\text{max}}$   
for each  $I_{ki}(x)$ , and we need  $N_{\text{max}}$  of them.
- We therefore need memory space for at least  
 $N_{\text{max}} (N_{\text{mesh}})^3 > 27 (N_{\text{max}})^4$   
numbers
- With single precision (Float32), already 27 GB for  $N_{\text{max}}=128$ .

# Naive parallelization of the estimator

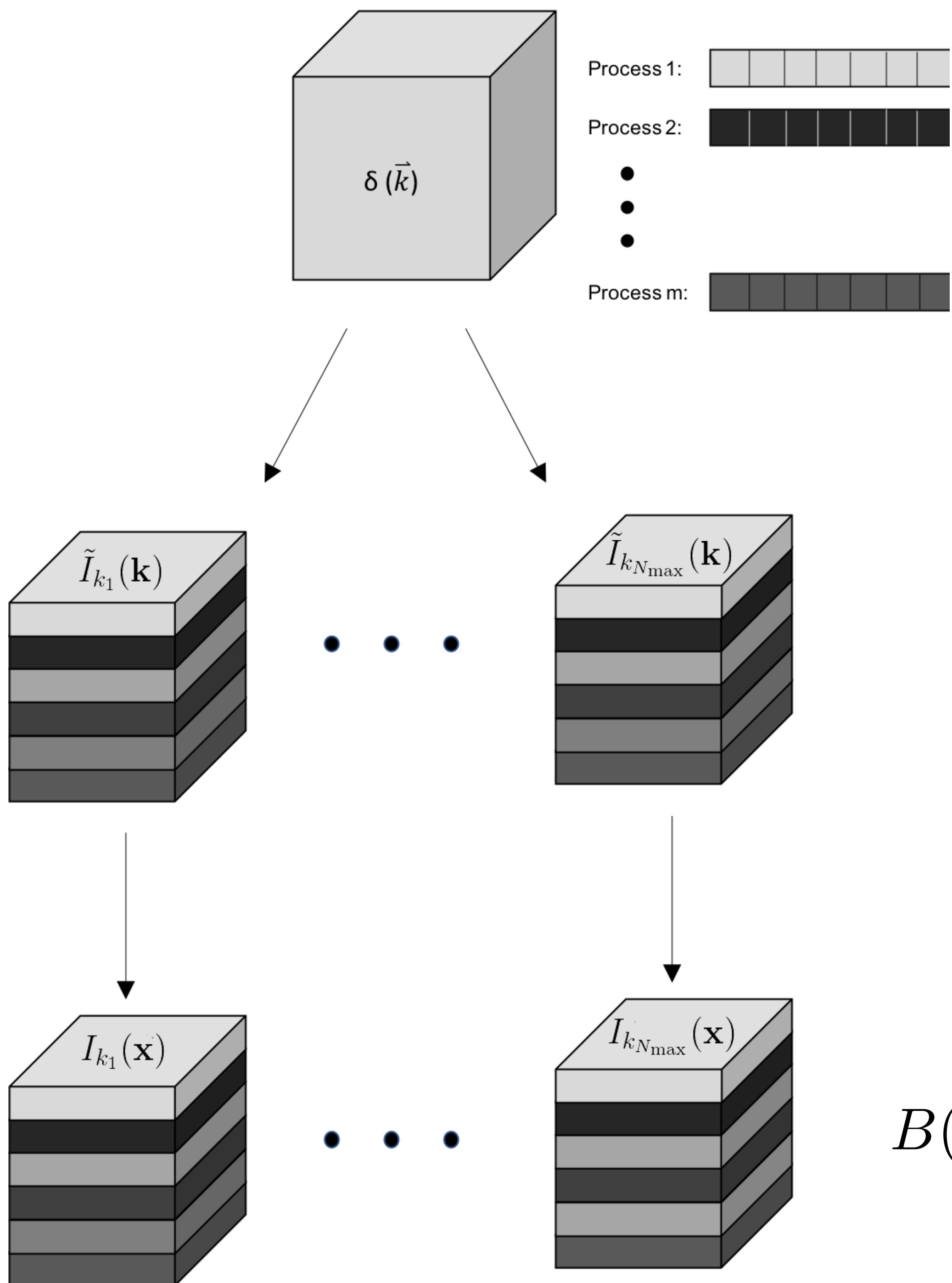
- Naive parallelization : Run each  $I_{k_i}(\mathbf{x})$  on one CPU

- Why bad?

We need to pull out the full 3D array to calculate the inner product

$$B(k_1, k_2, k_3) = \frac{V_f}{V_{(123)}^B (2\pi)^3} \int \frac{d^3x}{(2\pi)^3} I_{k_1}(\mathbf{x}) I_{k_2}(\mathbf{x}) I_{k_3}(\mathbf{x})$$

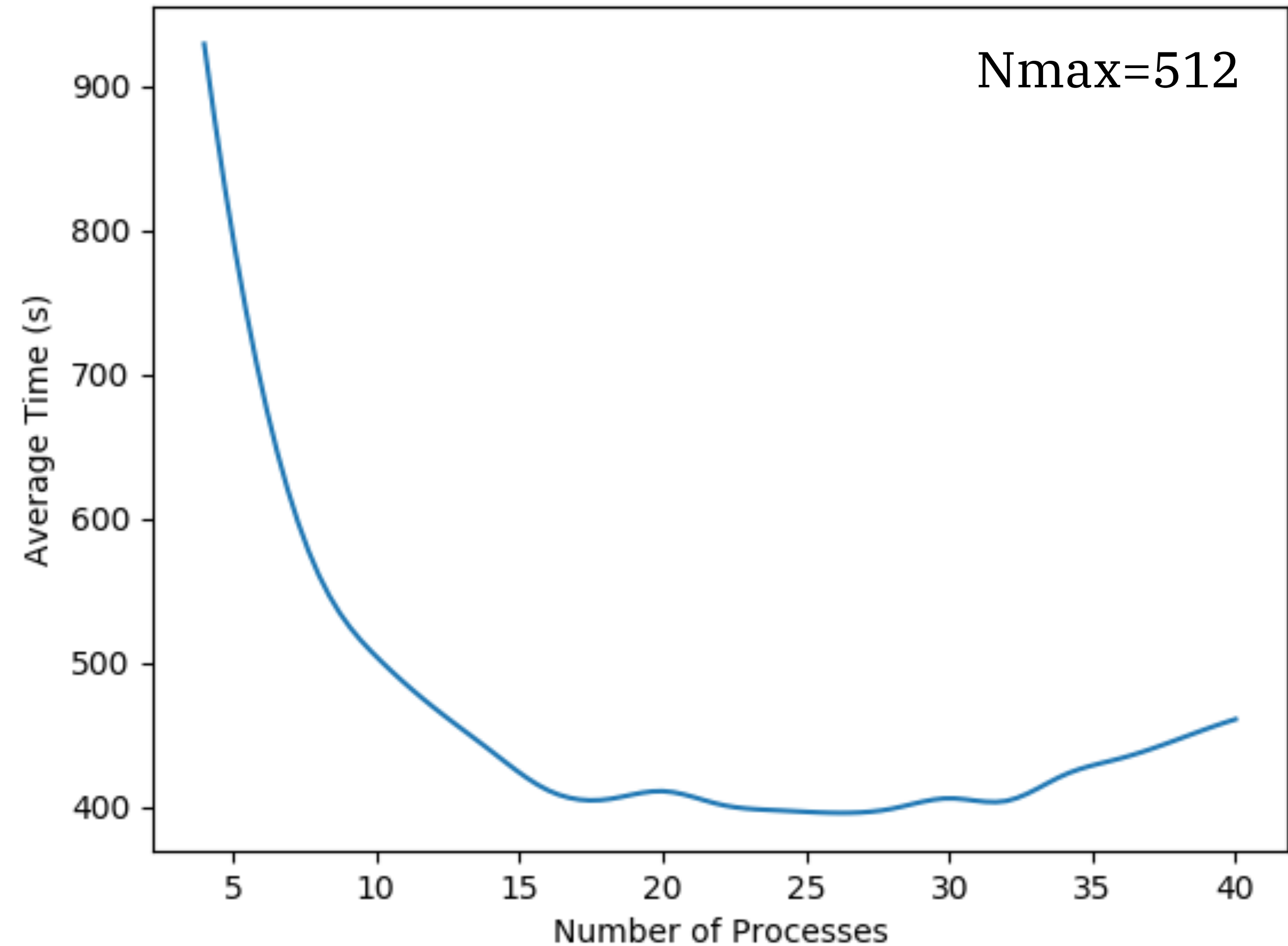
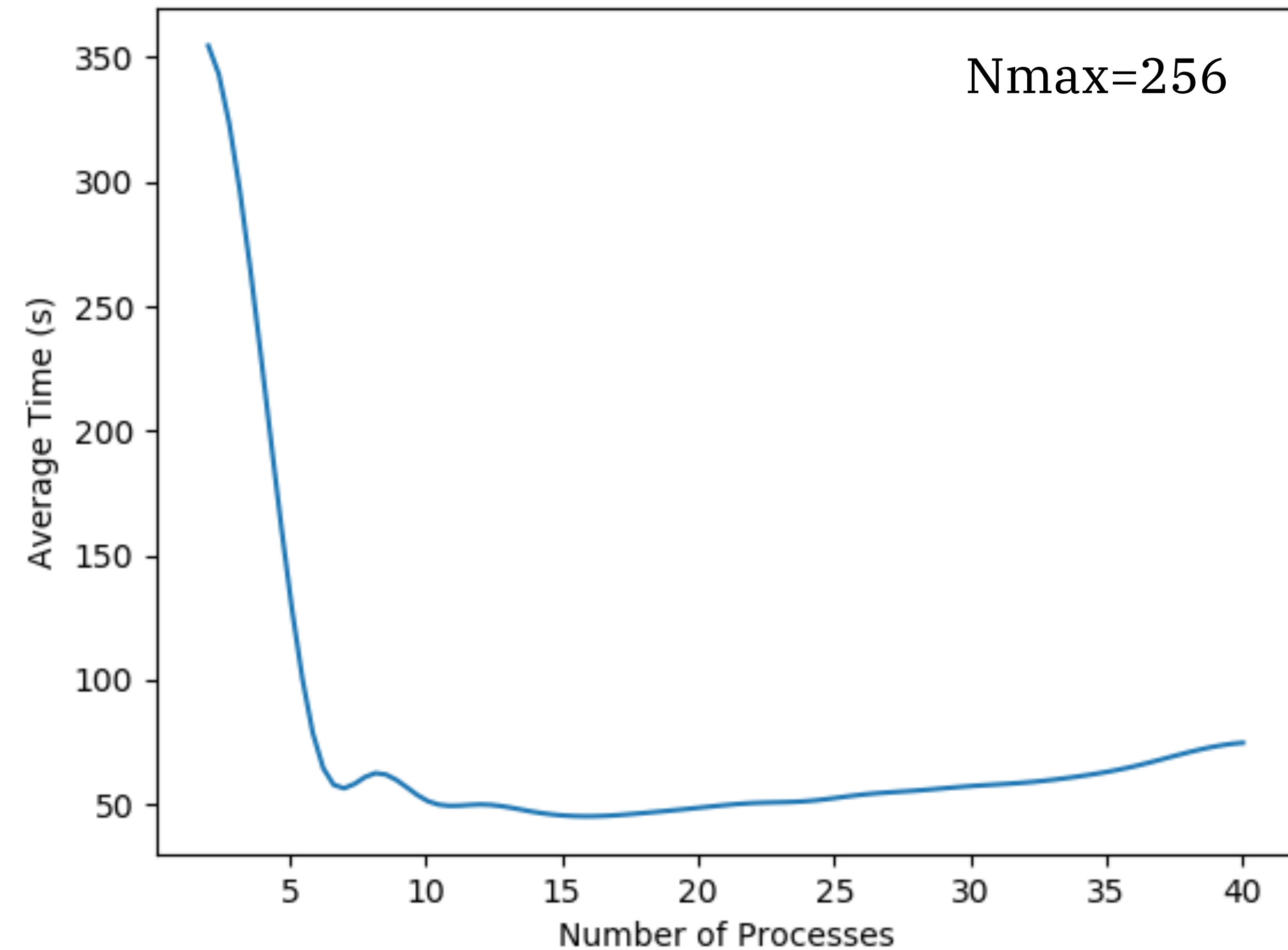
# Efficient parallelization



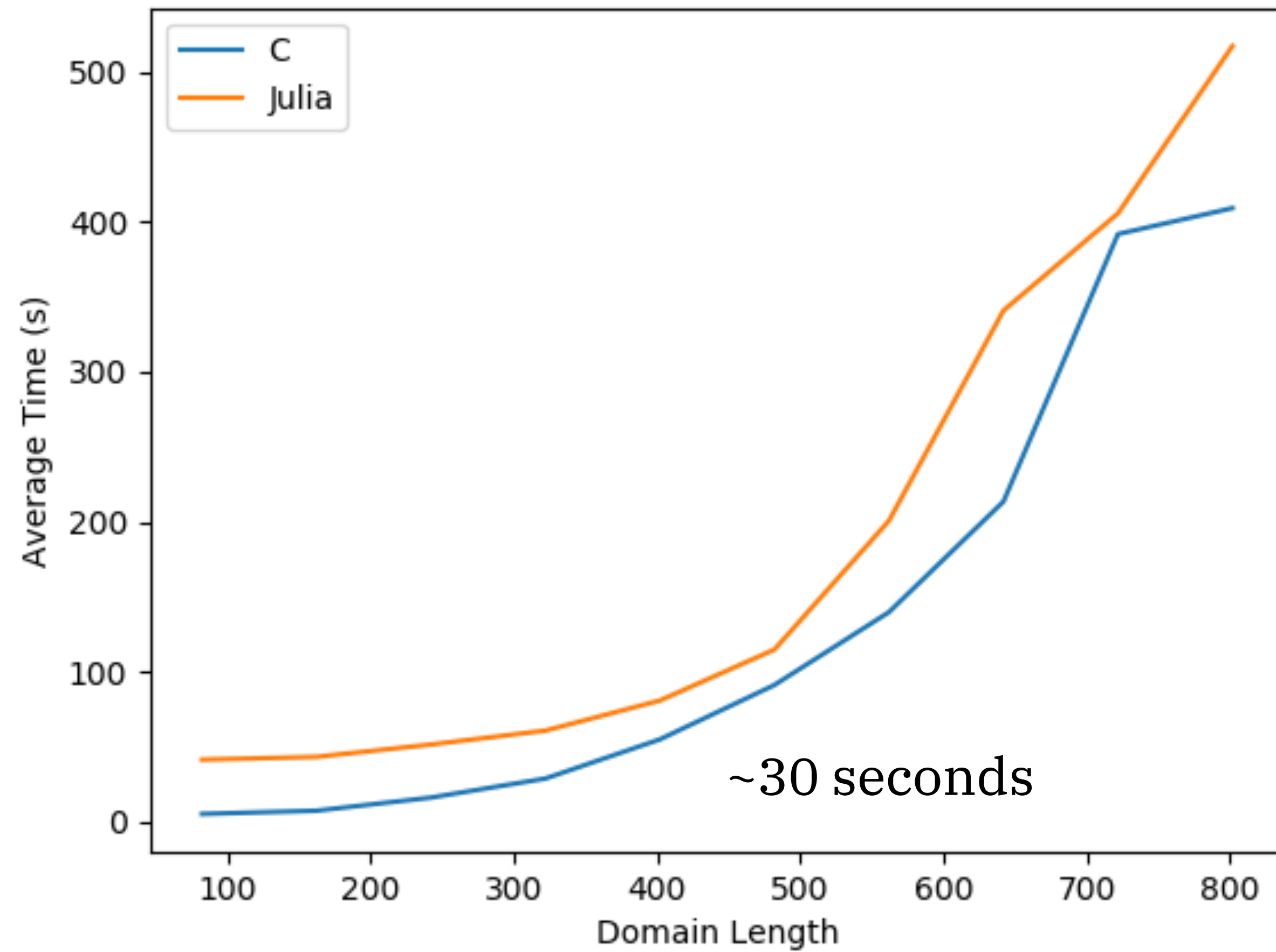
- Multiplication only done locally!
- Minimize the interCPU communication:
  - When FFT the last dimension
  - When reducing the sum

$$B(k_1, k_2, k_3) = \frac{V_f}{V_{(123)}^B (2\pi)^3} \int \frac{d^3 x}{(2\pi)^3} I_{k_1}(\mathbf{x}) I_{k_2}(\mathbf{x}) I_{k_3}(\mathbf{x})$$

# Efficient parallelization, result

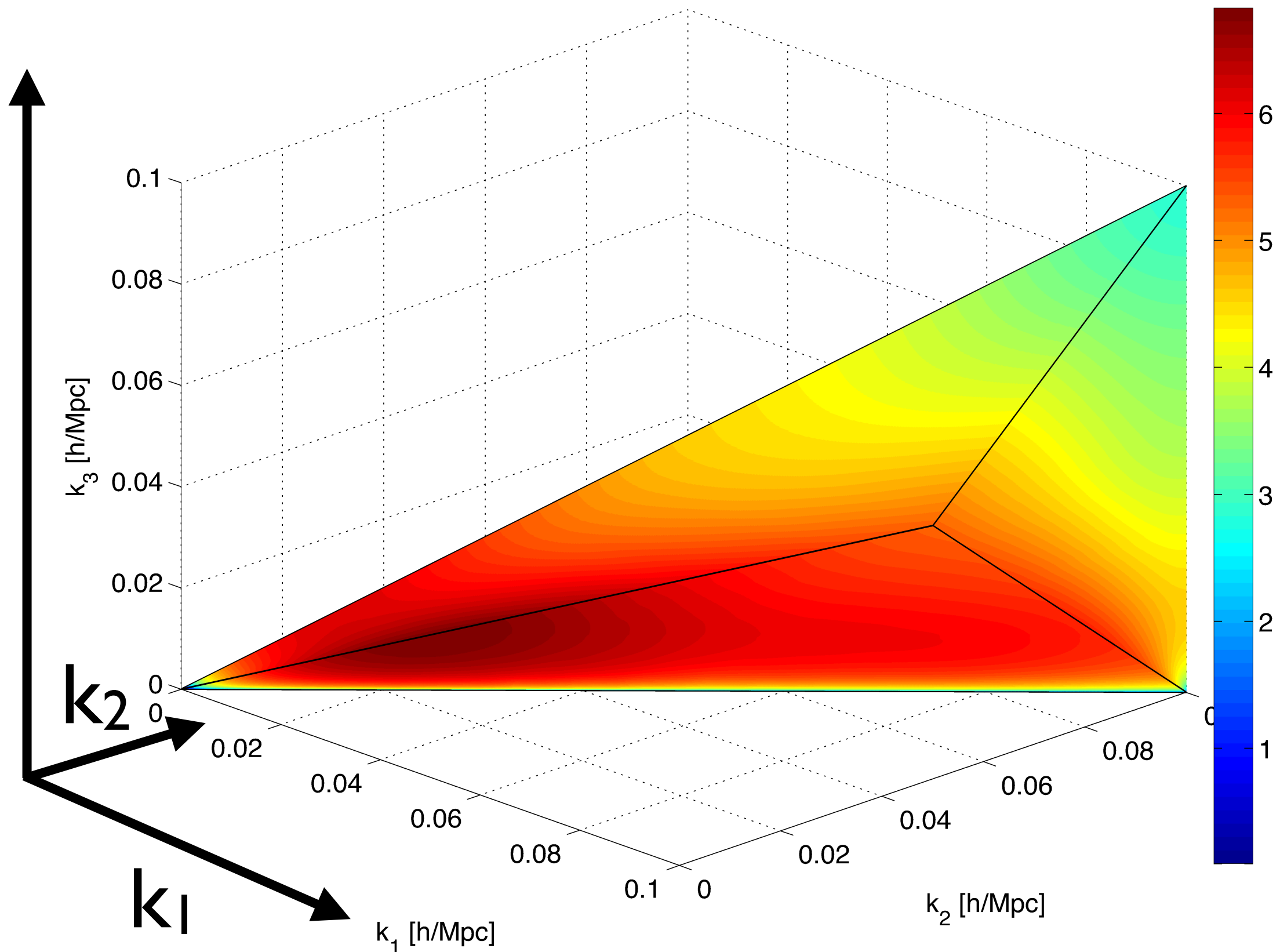
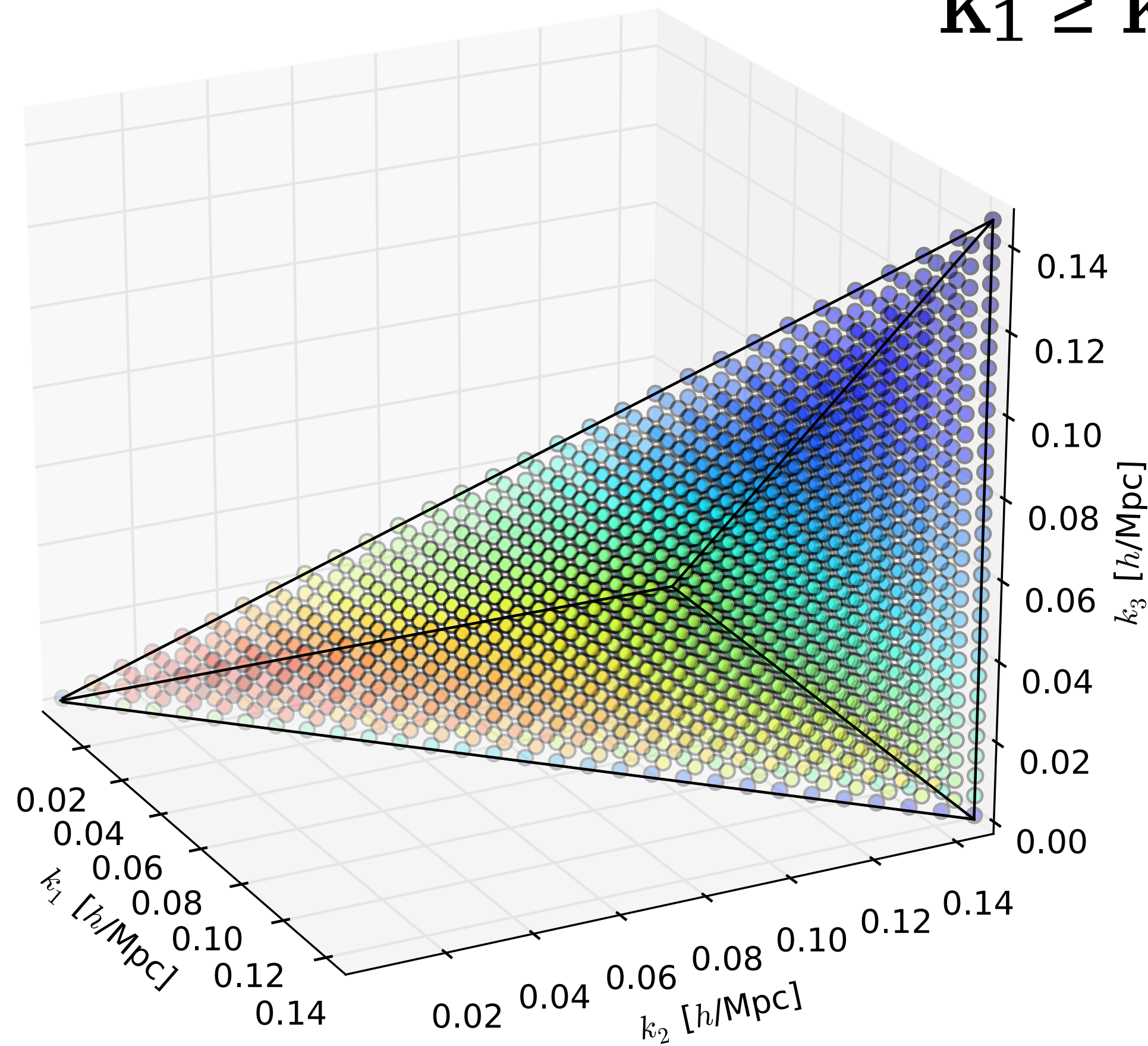


# C vs. Julia

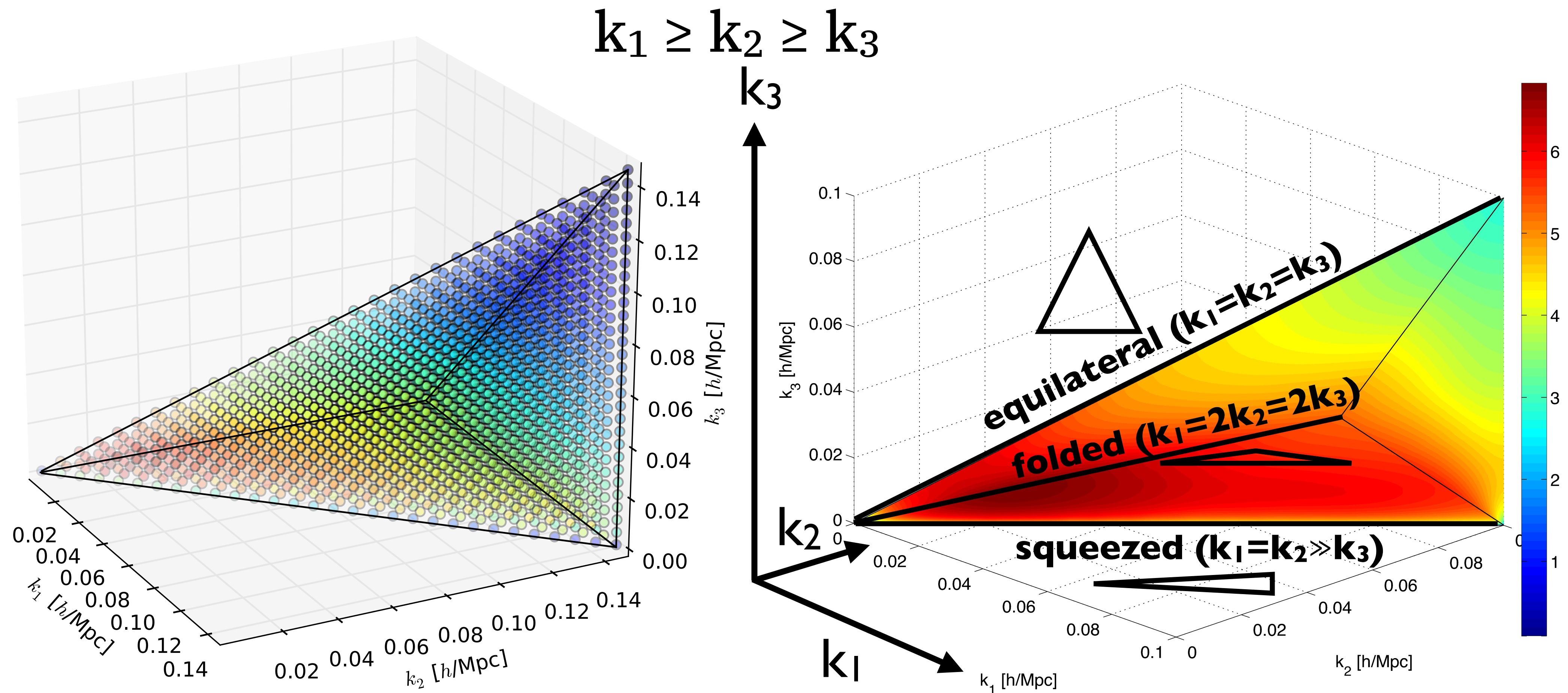


# Visualizing bispectrum

$$k_1 \geq k_2 \geq k_3$$

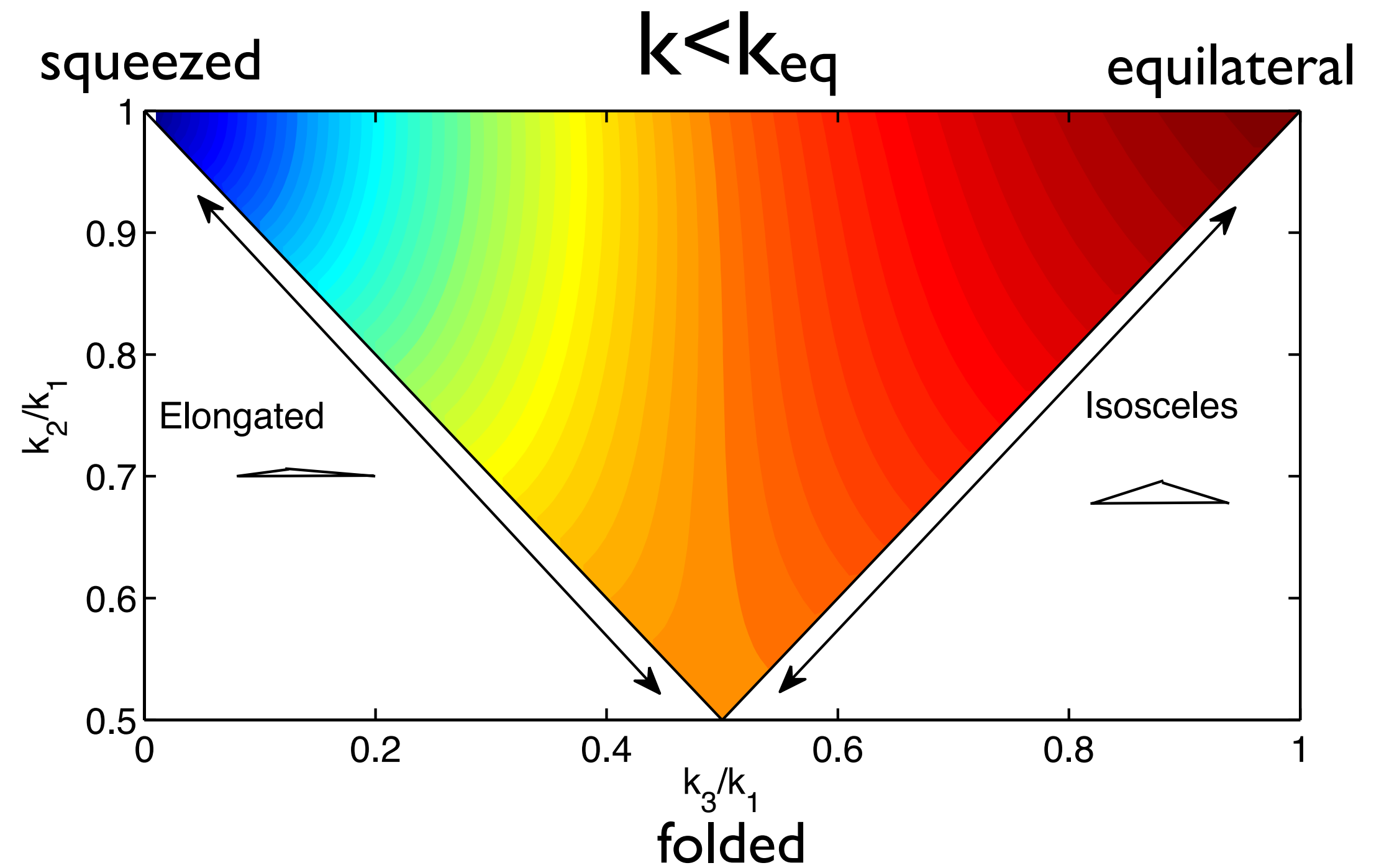
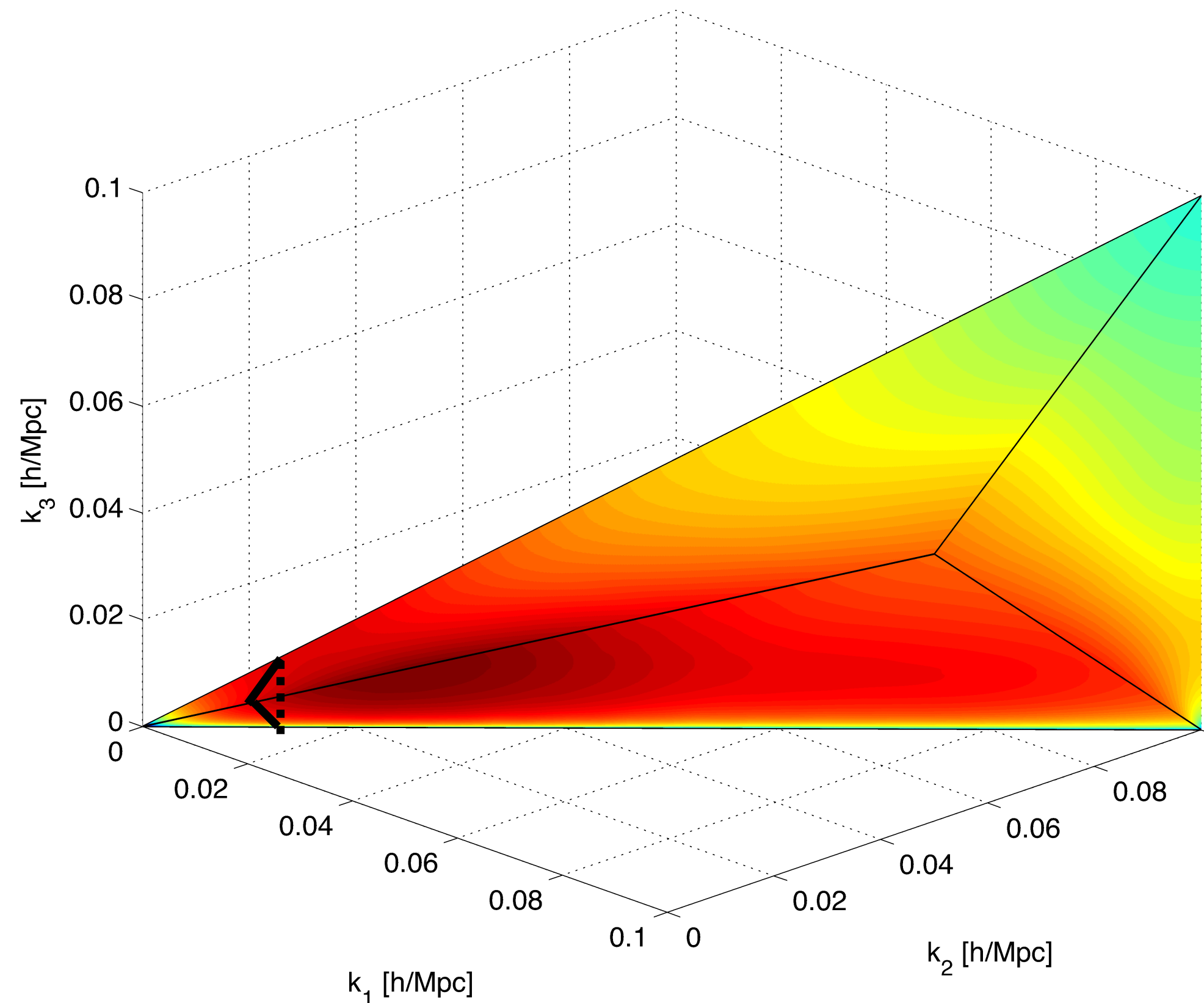


# Visualizing bispectrum



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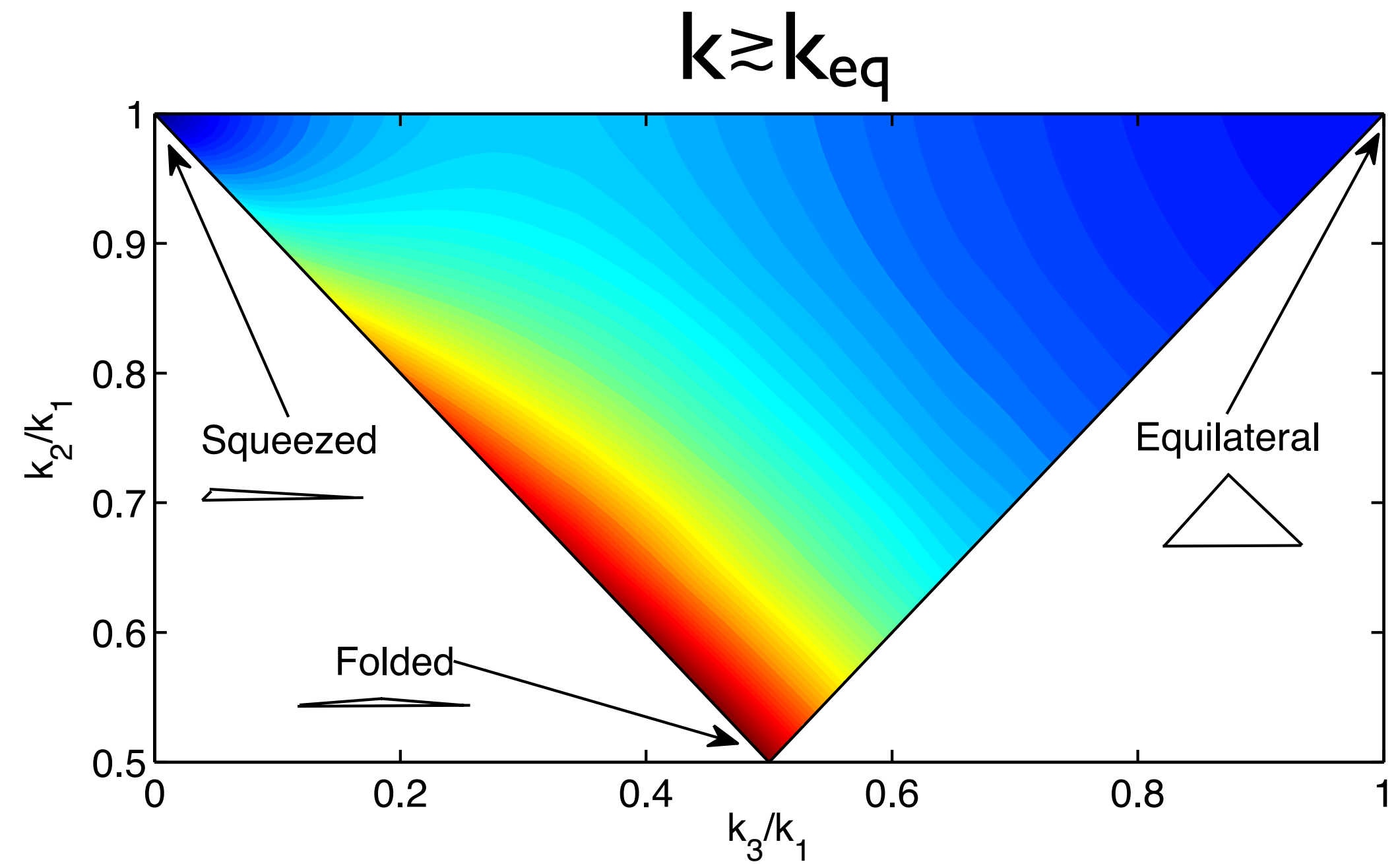
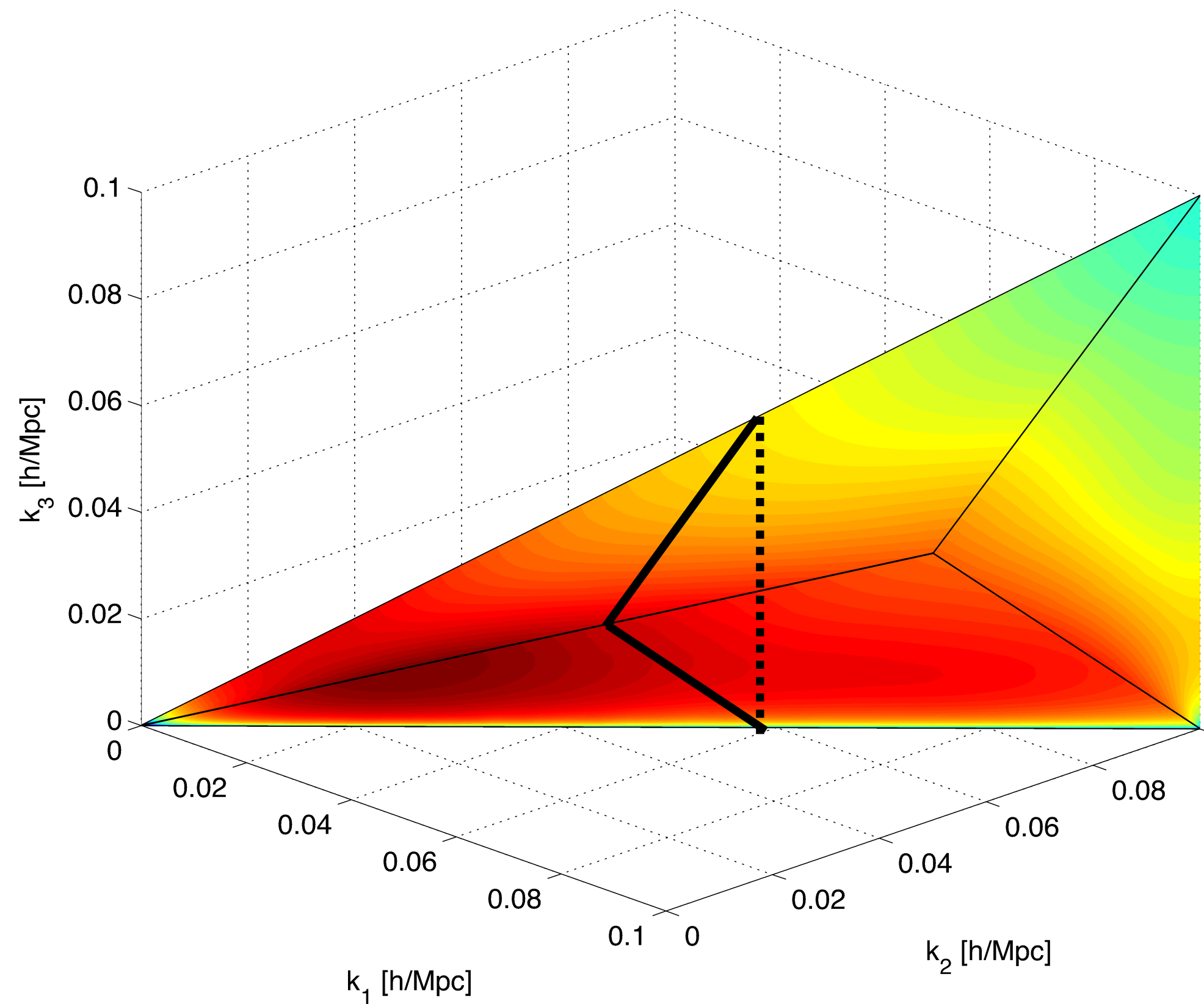


The slope of power spectrum  
 $d \ln P / d \ln k > 0$



# Visualizing bispectrum

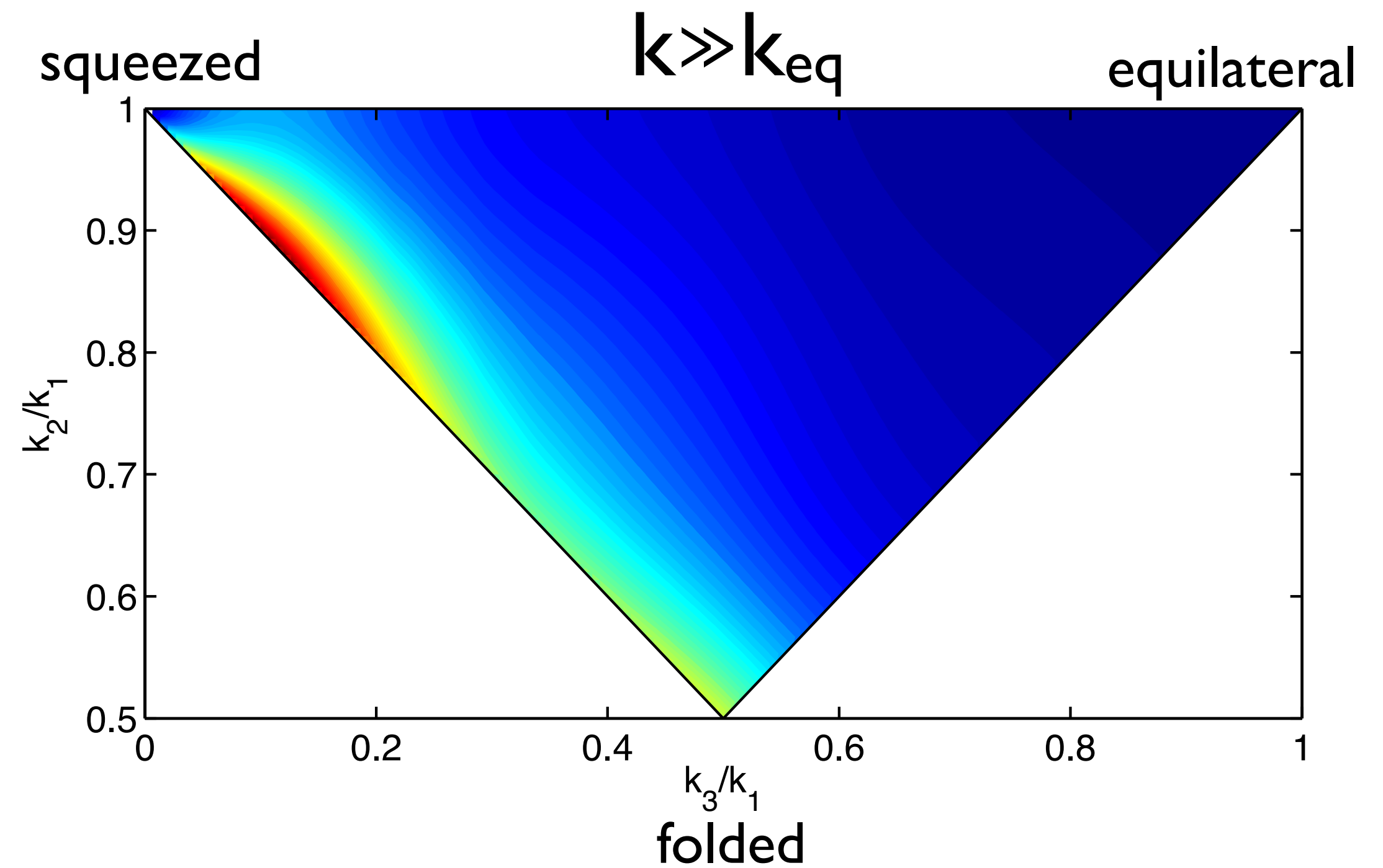
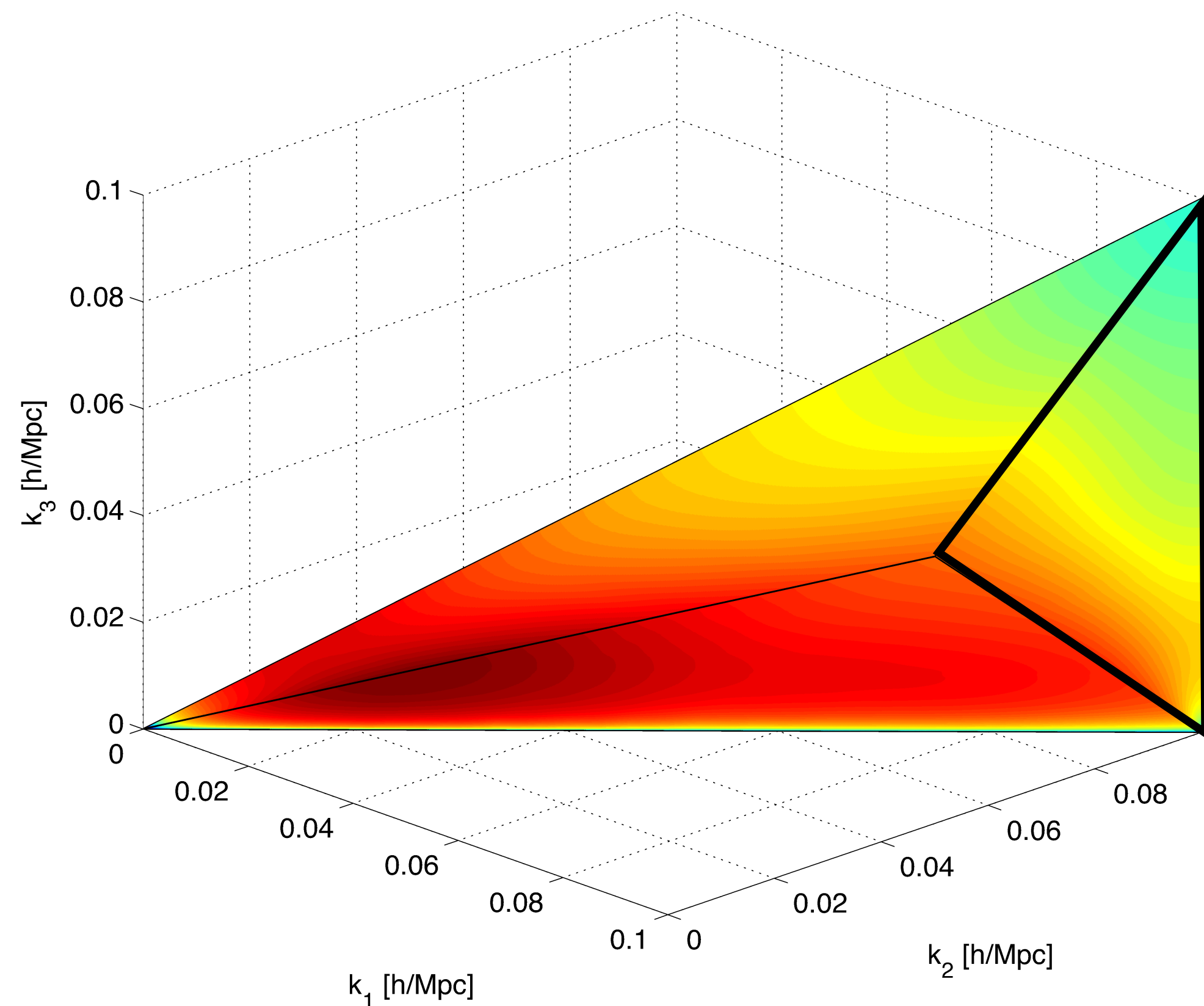
$$k_1 \geq k_2 \geq k_3$$



The slope of power spectrum  
 $d \ln P / d \ln k \approx 0$

# Visualizing bispectrum

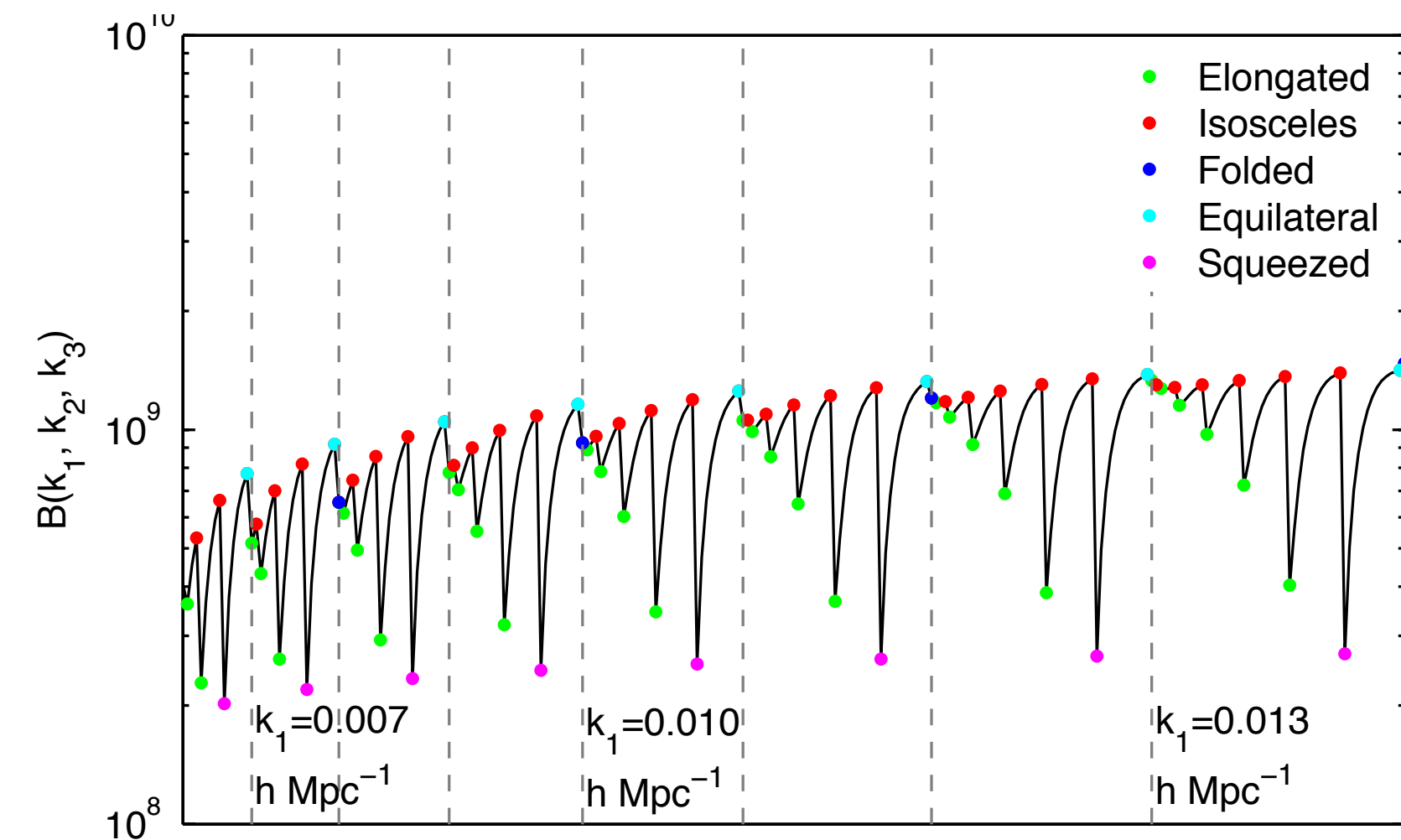
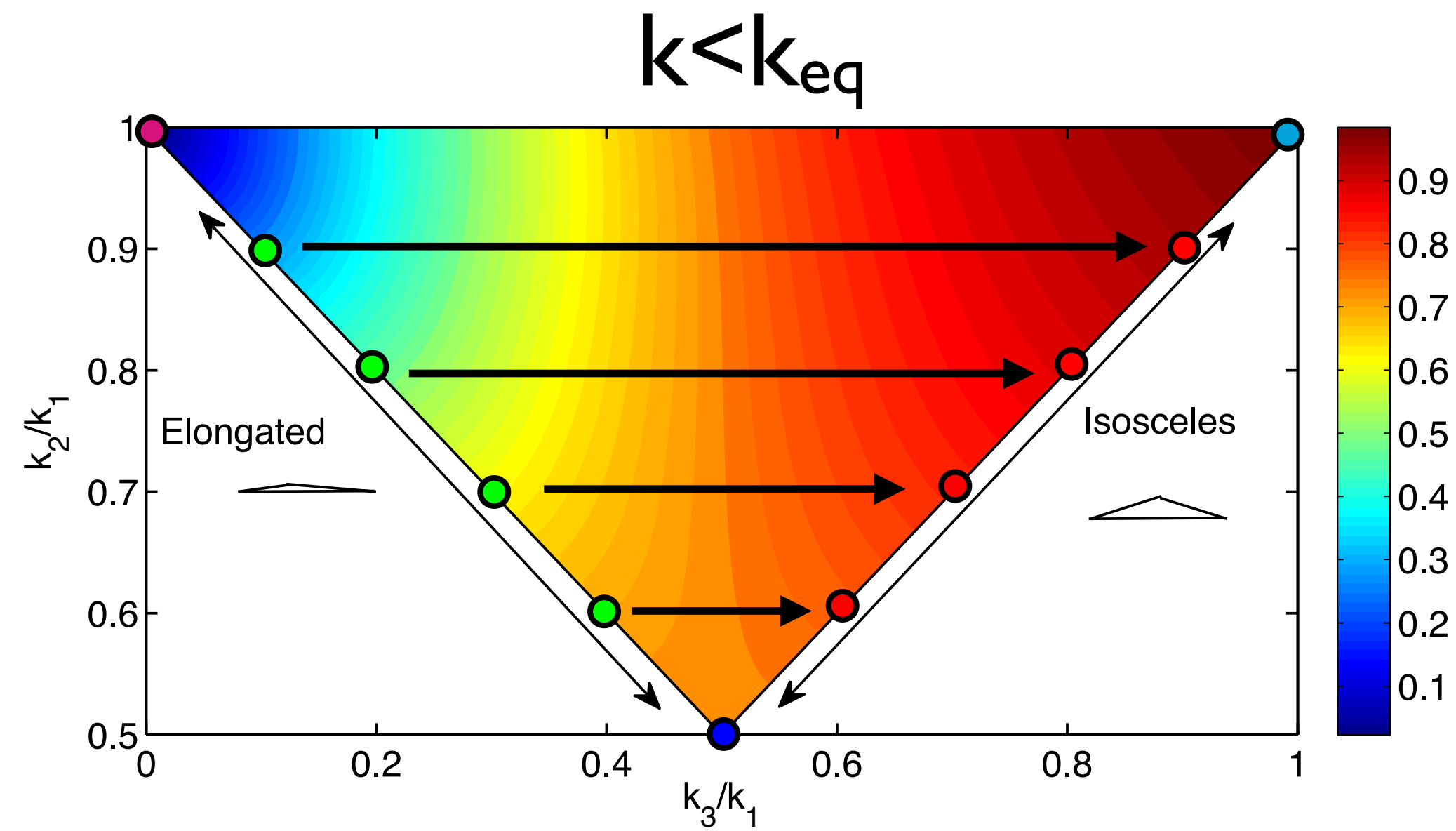
$$k_1 \geq k_2 \geq k_3$$



The slope of power spectrum  
 $d \ln P / d \ln k < 0$

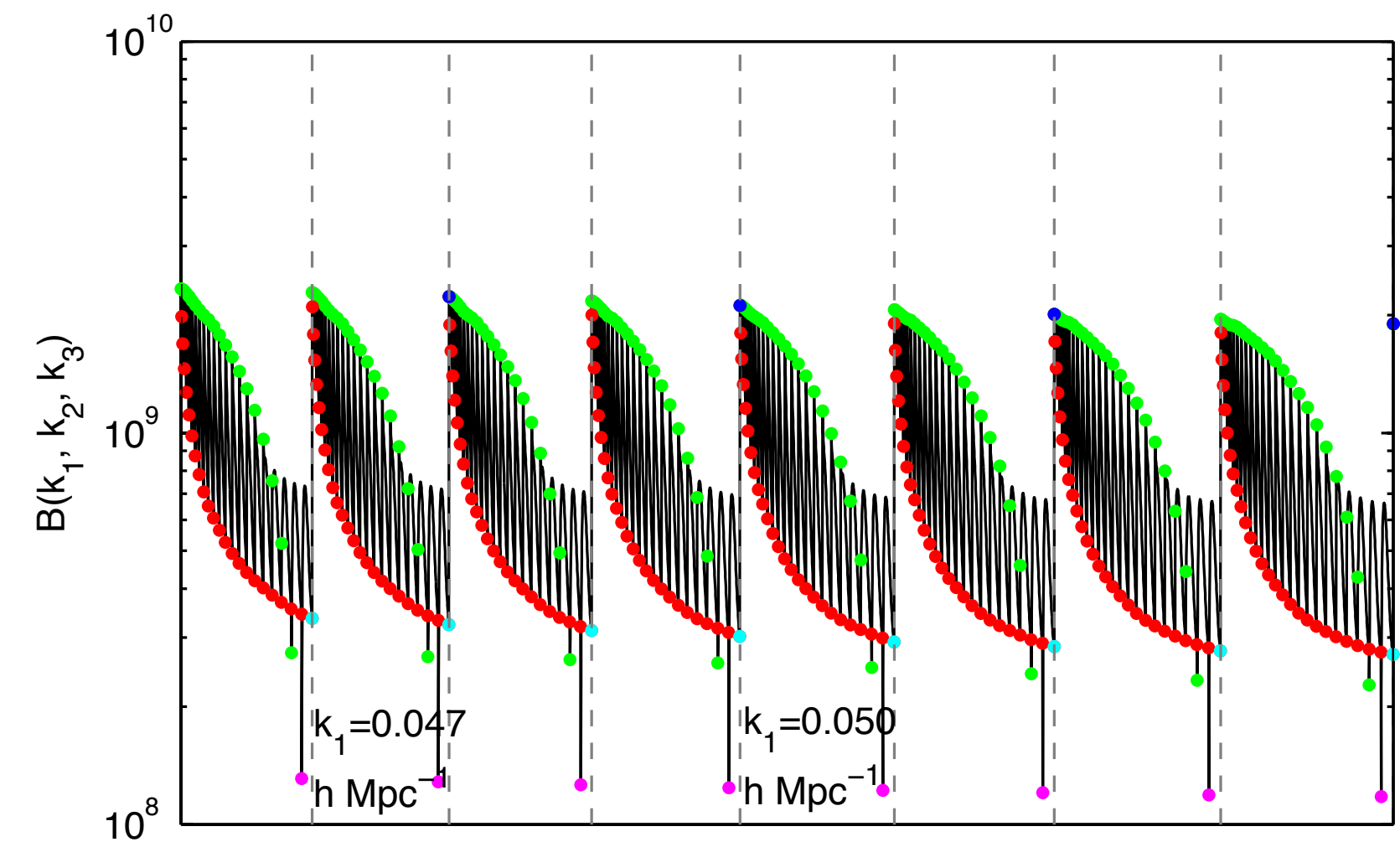
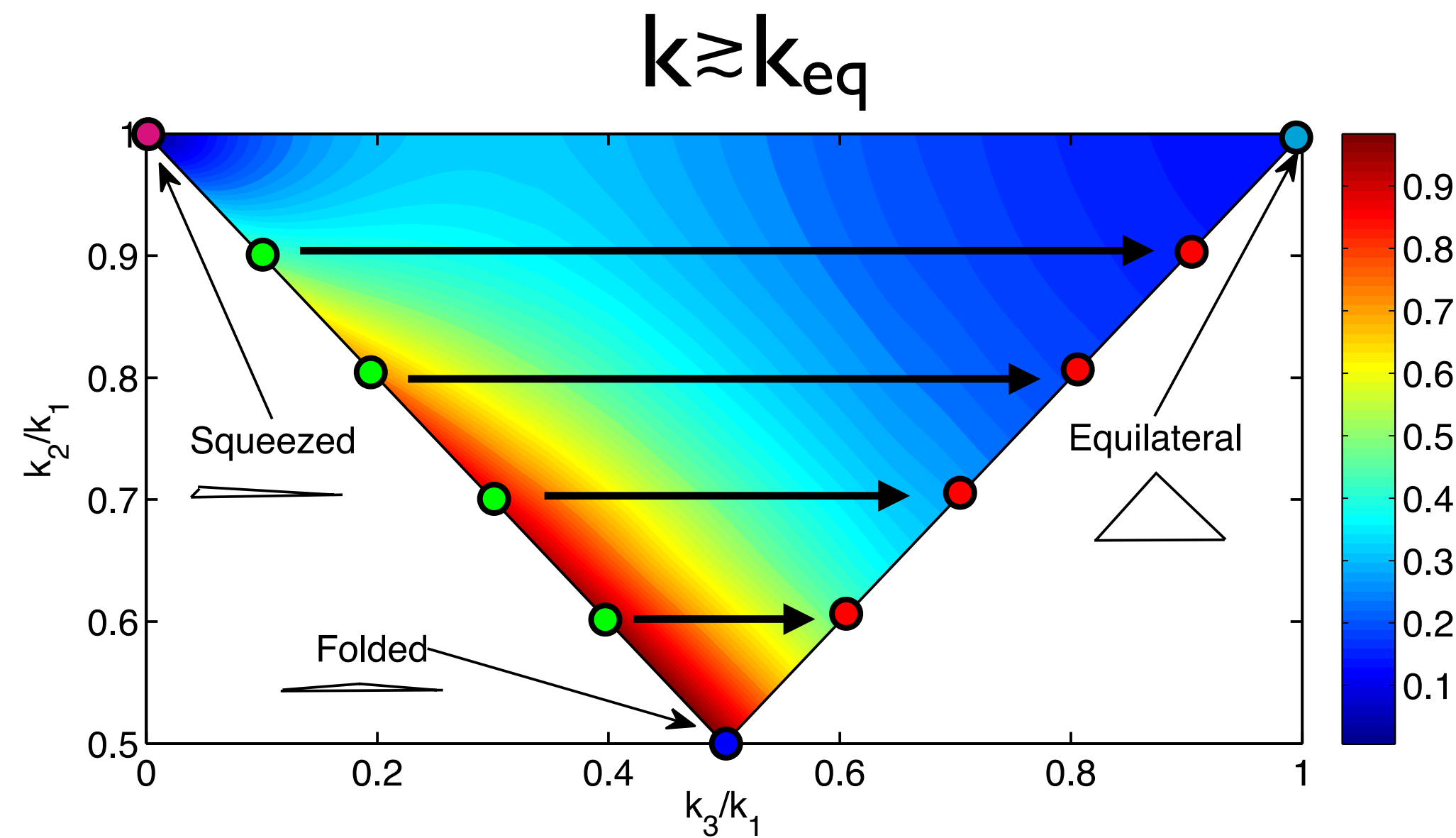
# Visualizing bispectrum

$$k_1 \geq k_2 \geq k_3$$



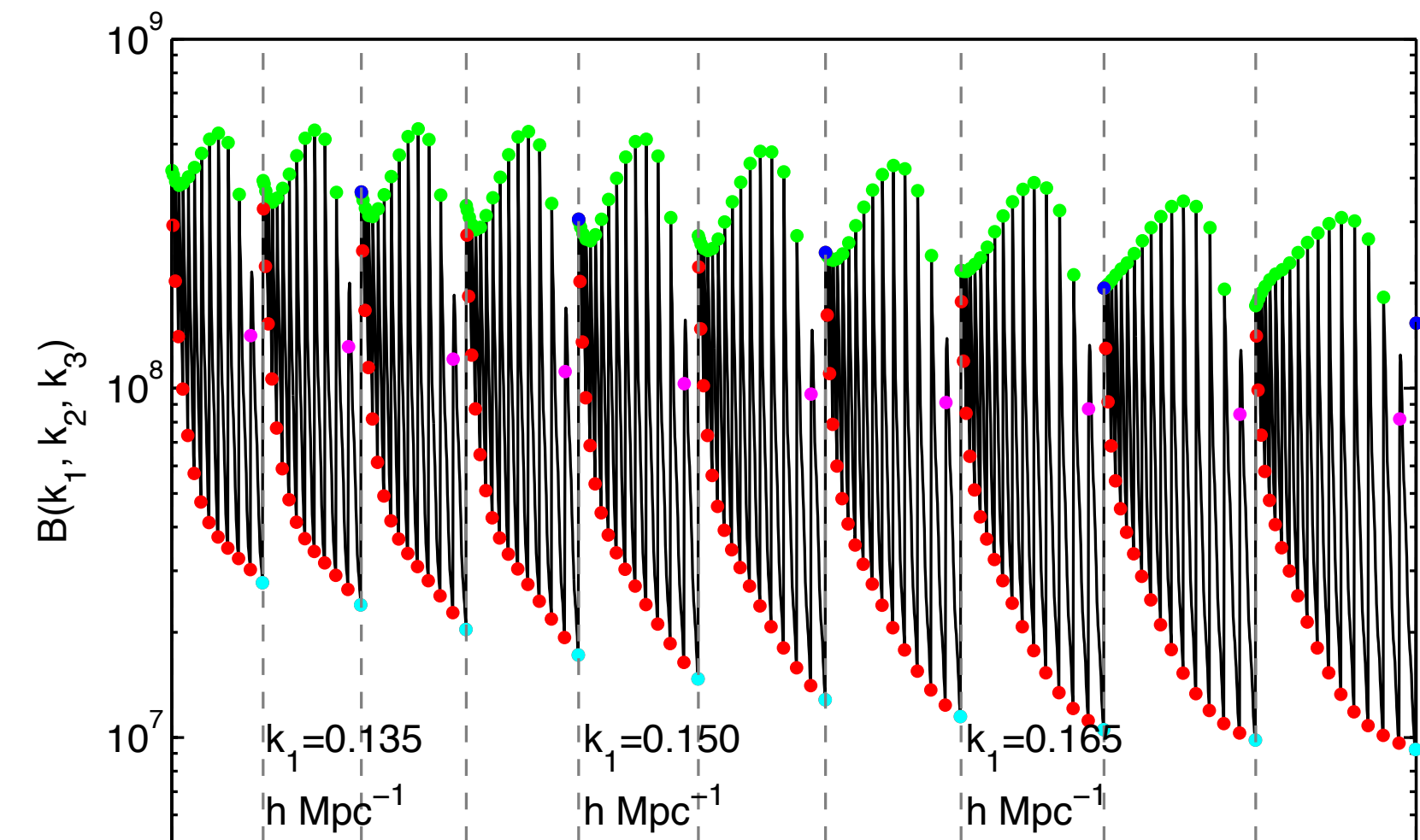
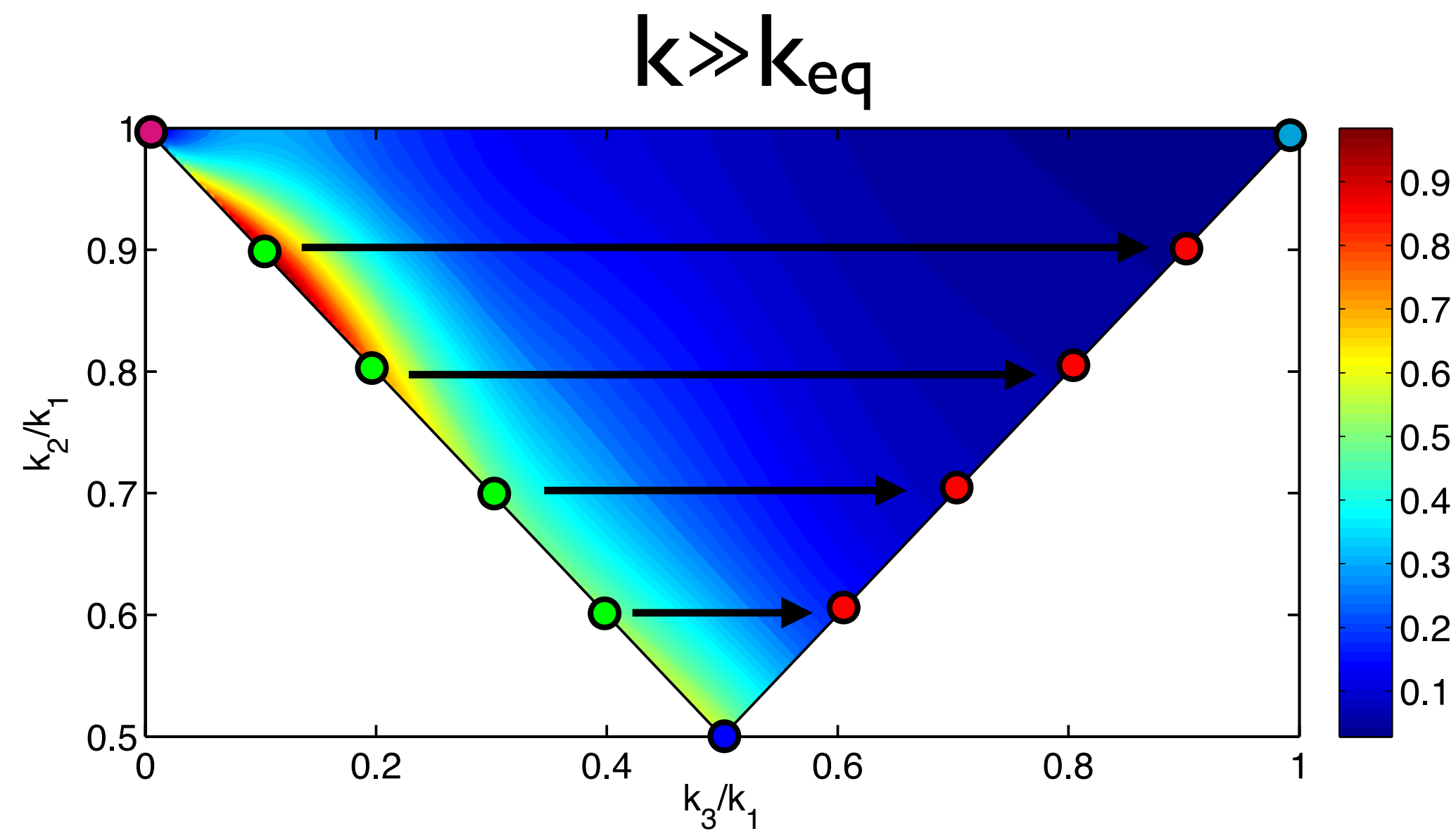
# Visualizing bispectrum

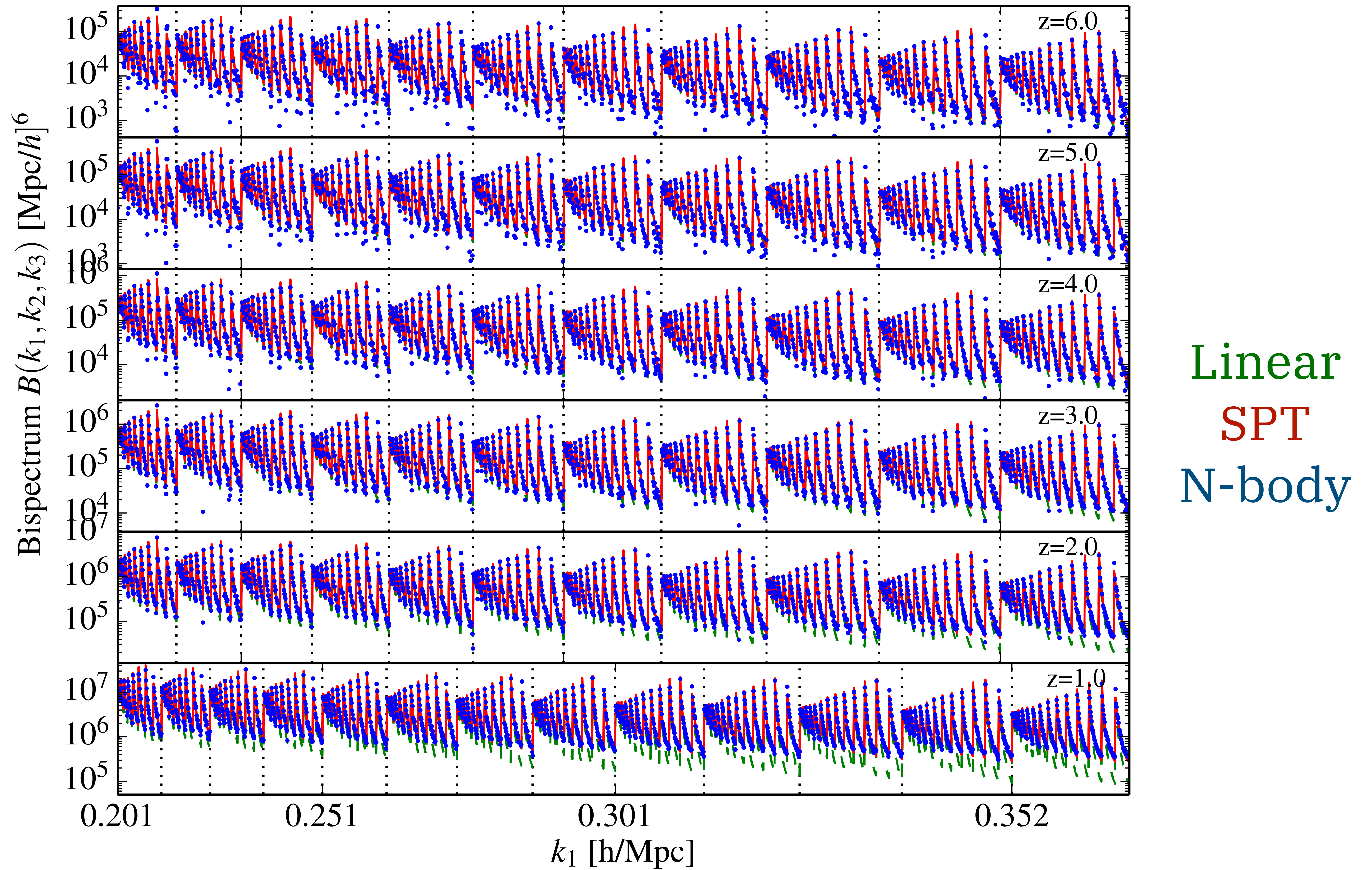
$$k_1 \geq k_2 \geq k_3$$



# Visualizing bispectrum

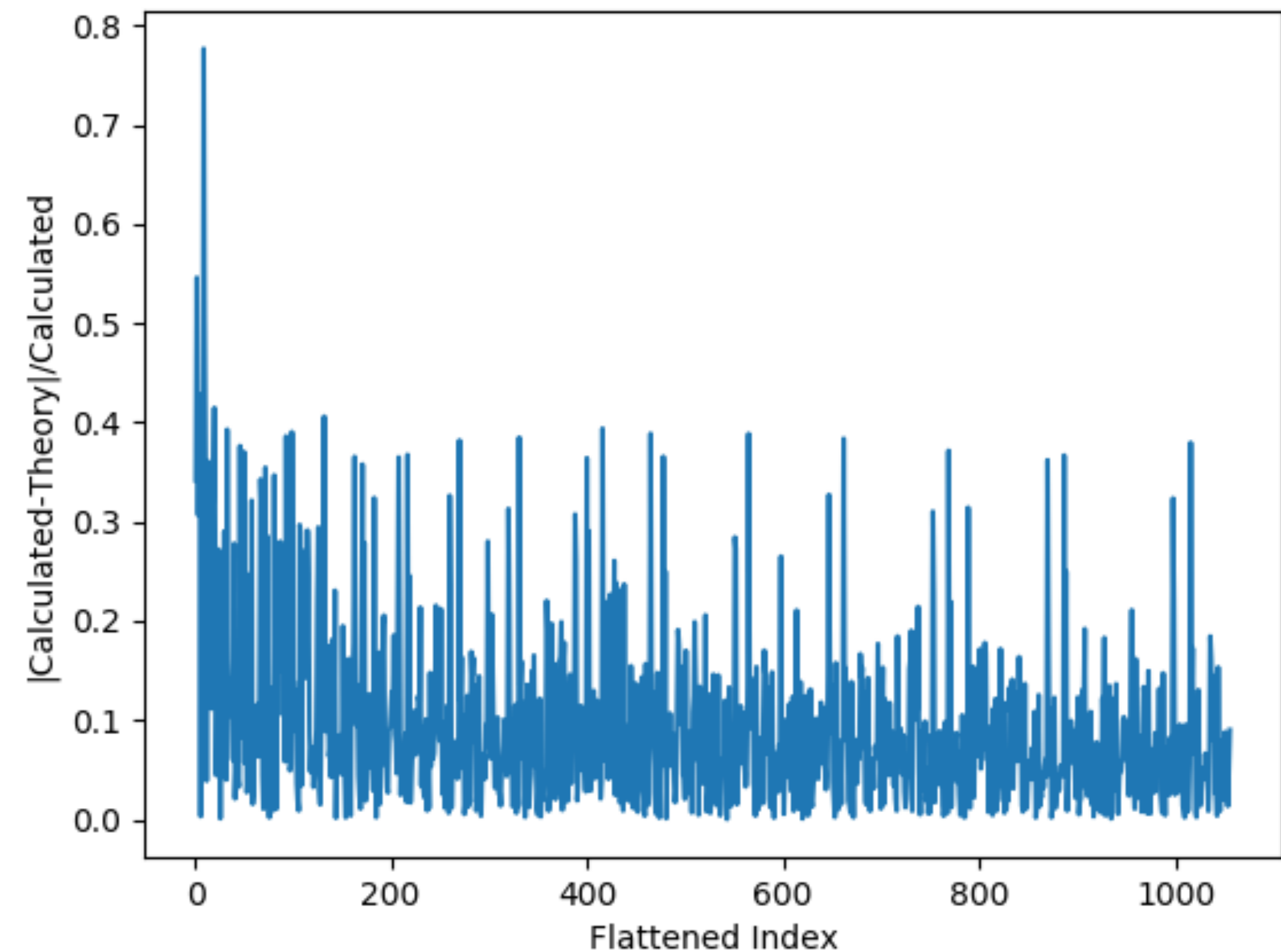
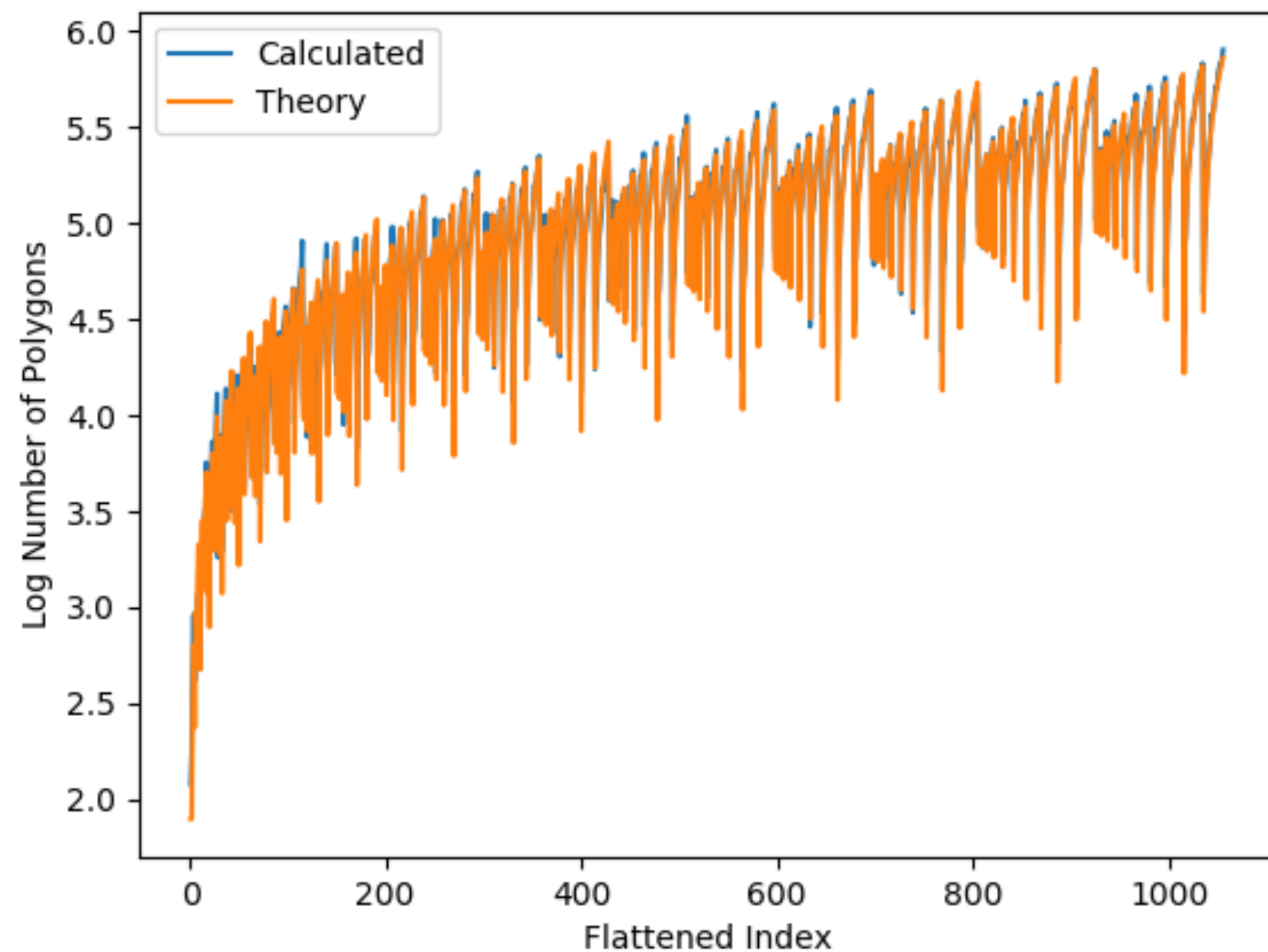
$$k_1 \geq k_2 \geq k_3$$





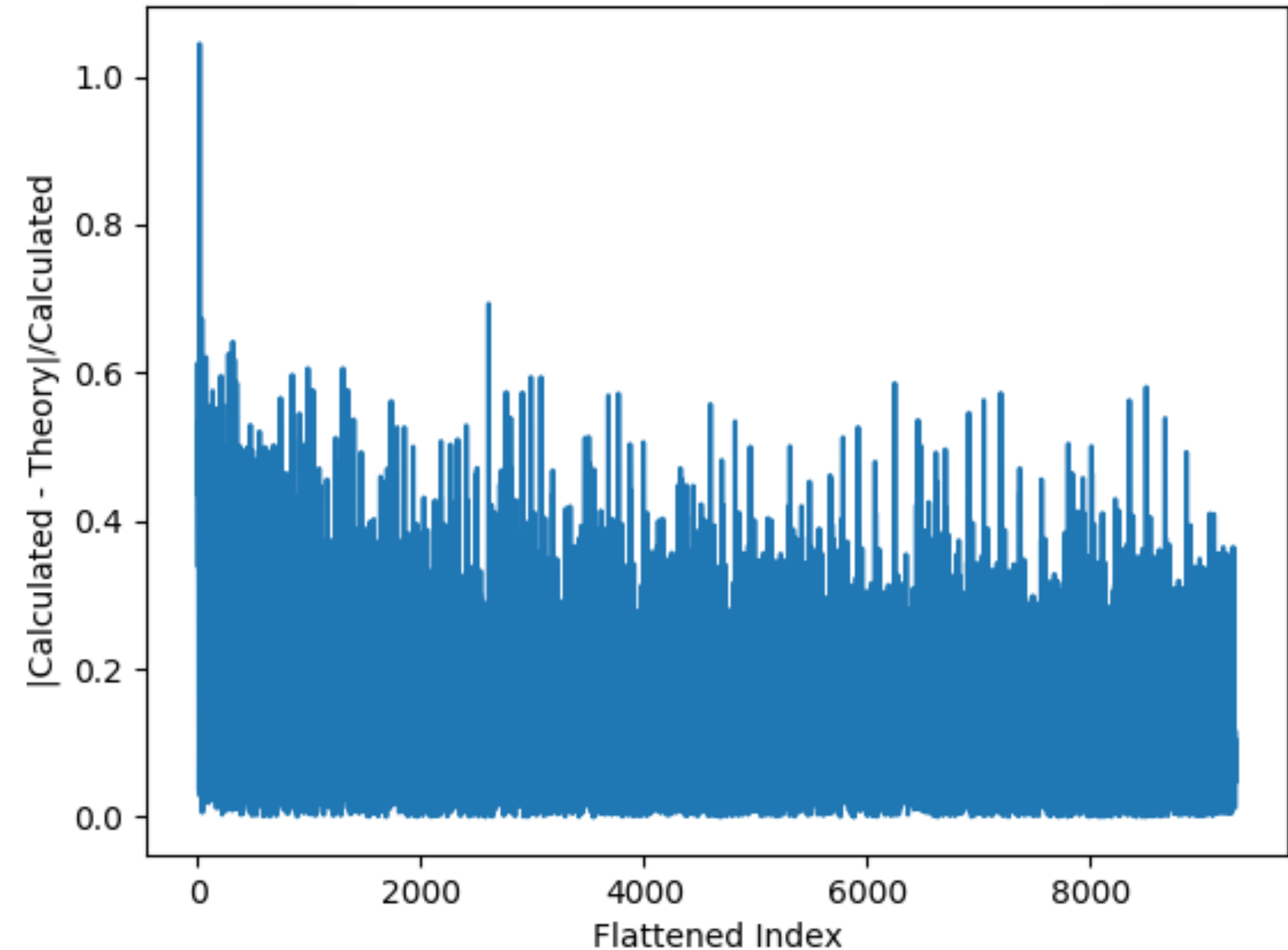
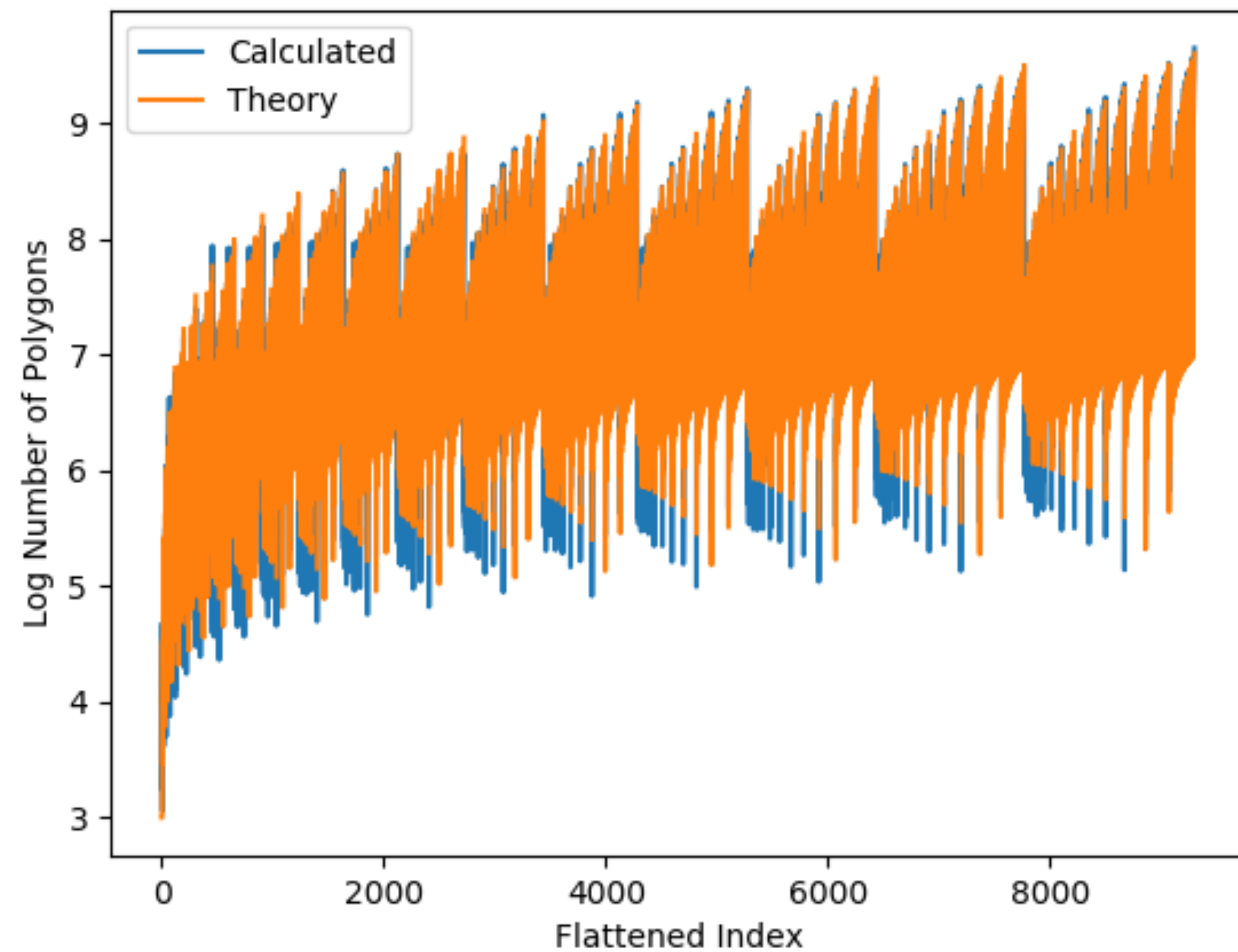
# Number of triangles

- Using the same estimator, but with  $\delta=1$ , we can calculate the total number of triangles. Spikes at  $k_1 = k_2 + k_3$



# Number of quadrilaterals

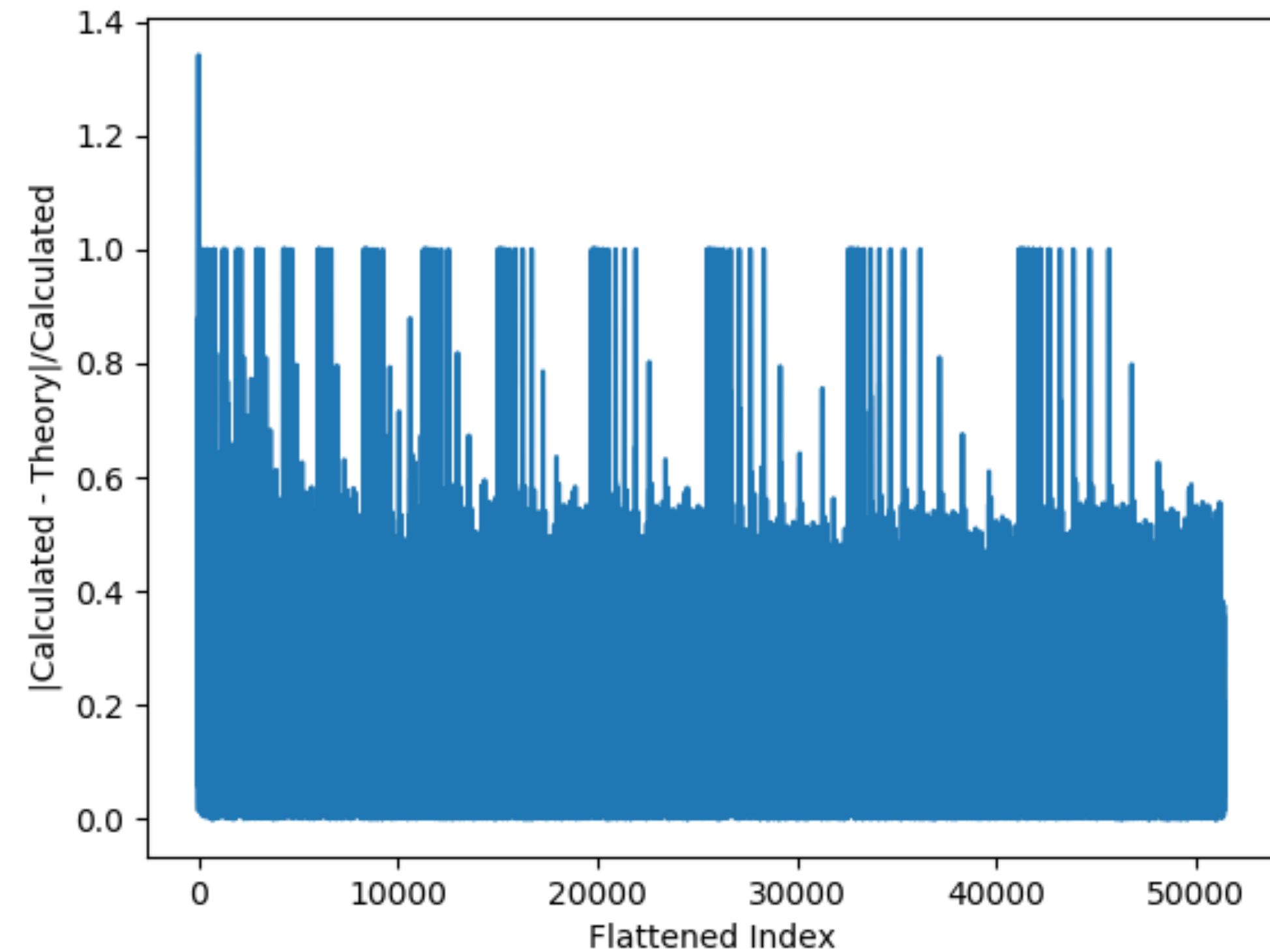
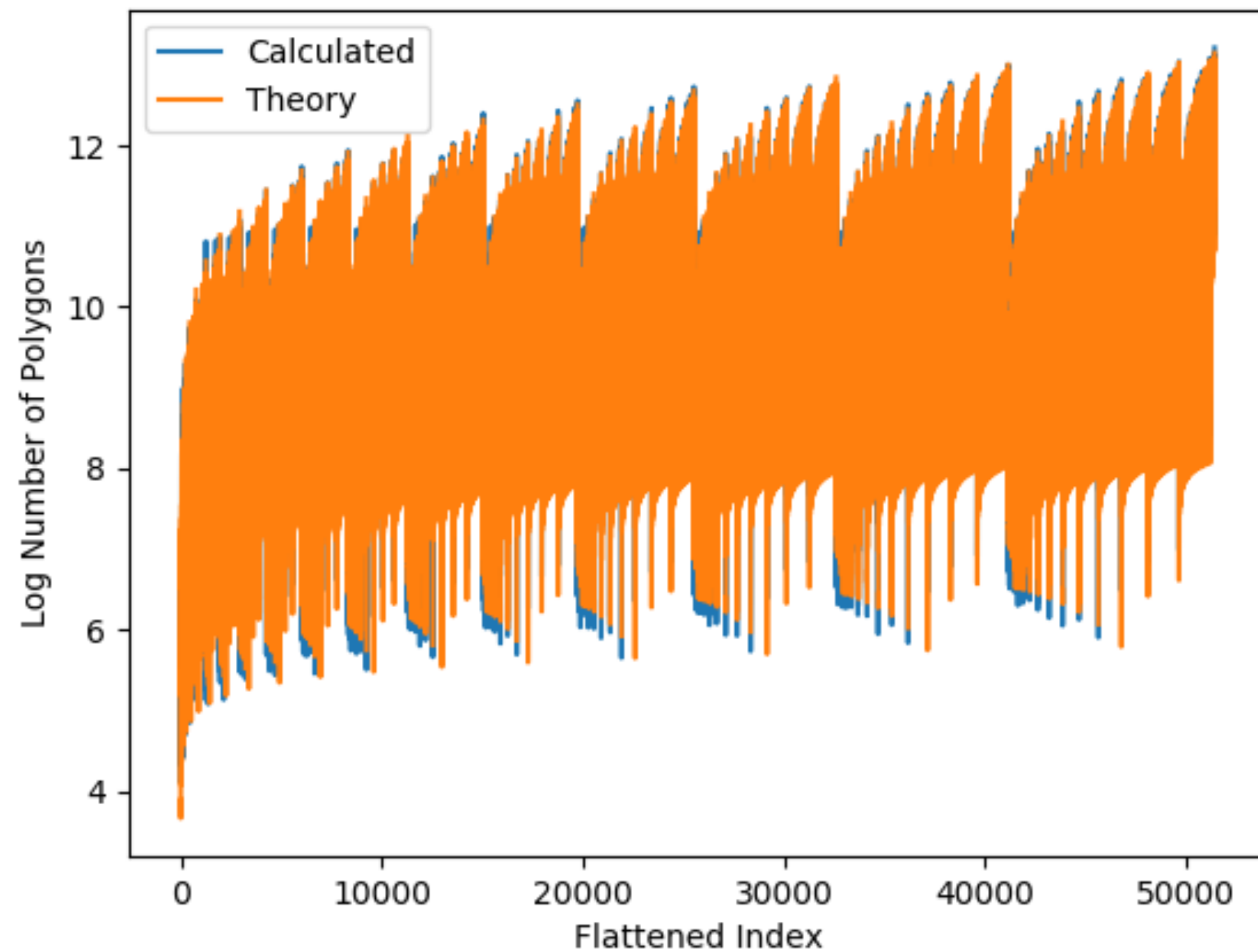
Why stop at triangle? Here's the angle averaged Trispectrum!





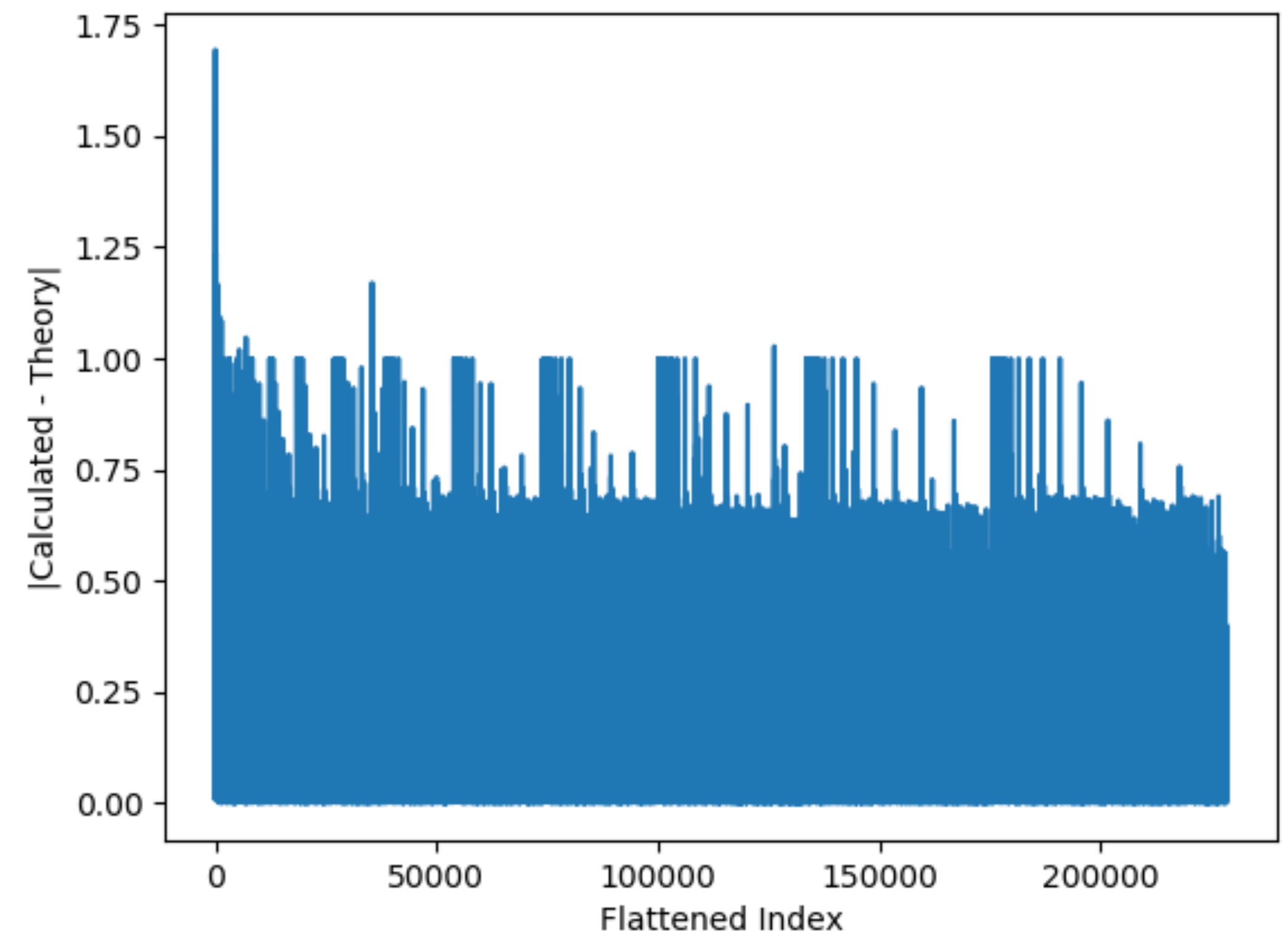
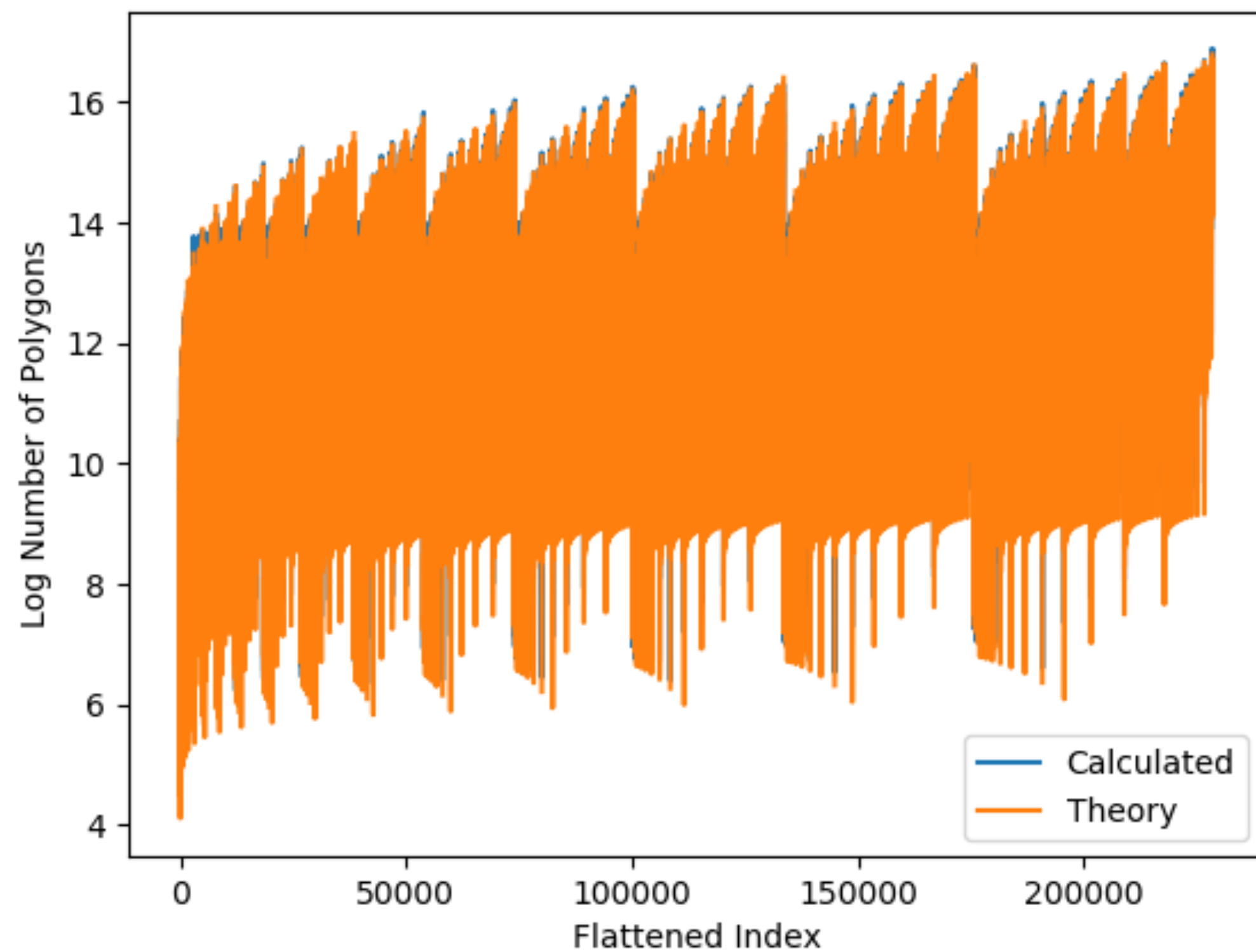
# Number of pentagons

Why stop at trispectrum? Here's the angle averaged quadspectrum!

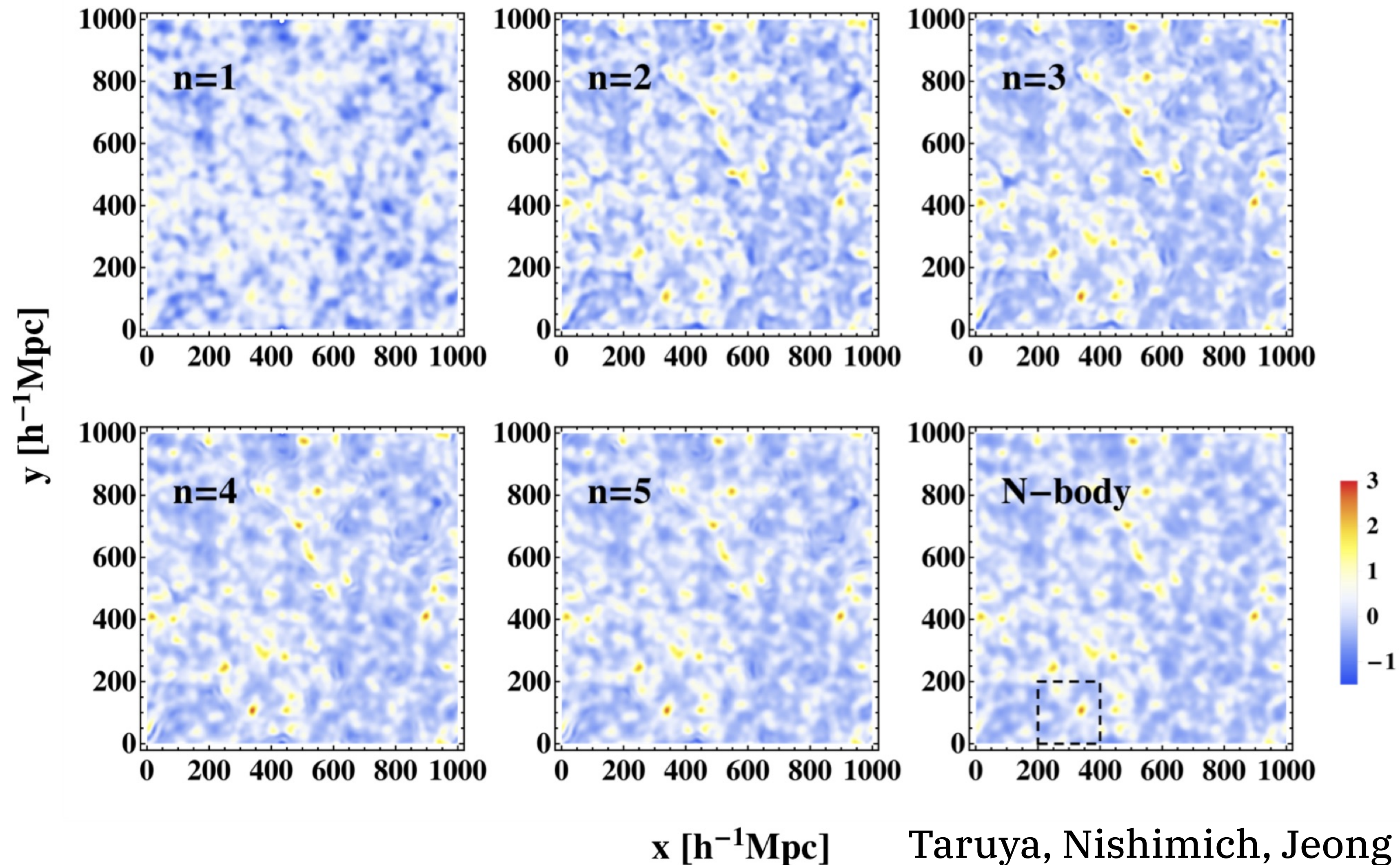


# Number of hexagons

Why stop at quadspectrum? Here's the angle averaged pentaspectrum!



# Application: polyspectra with GridSPT



Taruya, Nishimich, Jeong (2018)

# Conclusion

- We present an efficient parallel algorithm for calculating higher-order polyspectra
- With the parallelization, we can overcome the high memory requirement of Scoccimarro estimator, and the parallel version is quite fast!
- Applying it to GridSPT, we can calculate the SPT prediction for the higher-order polyspectra!